

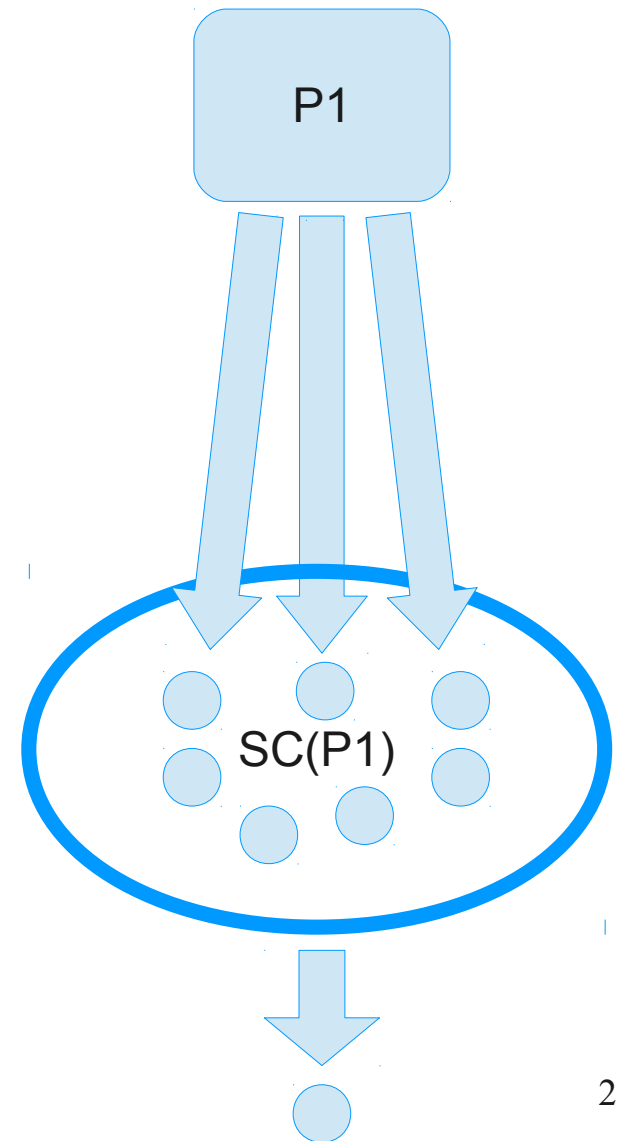
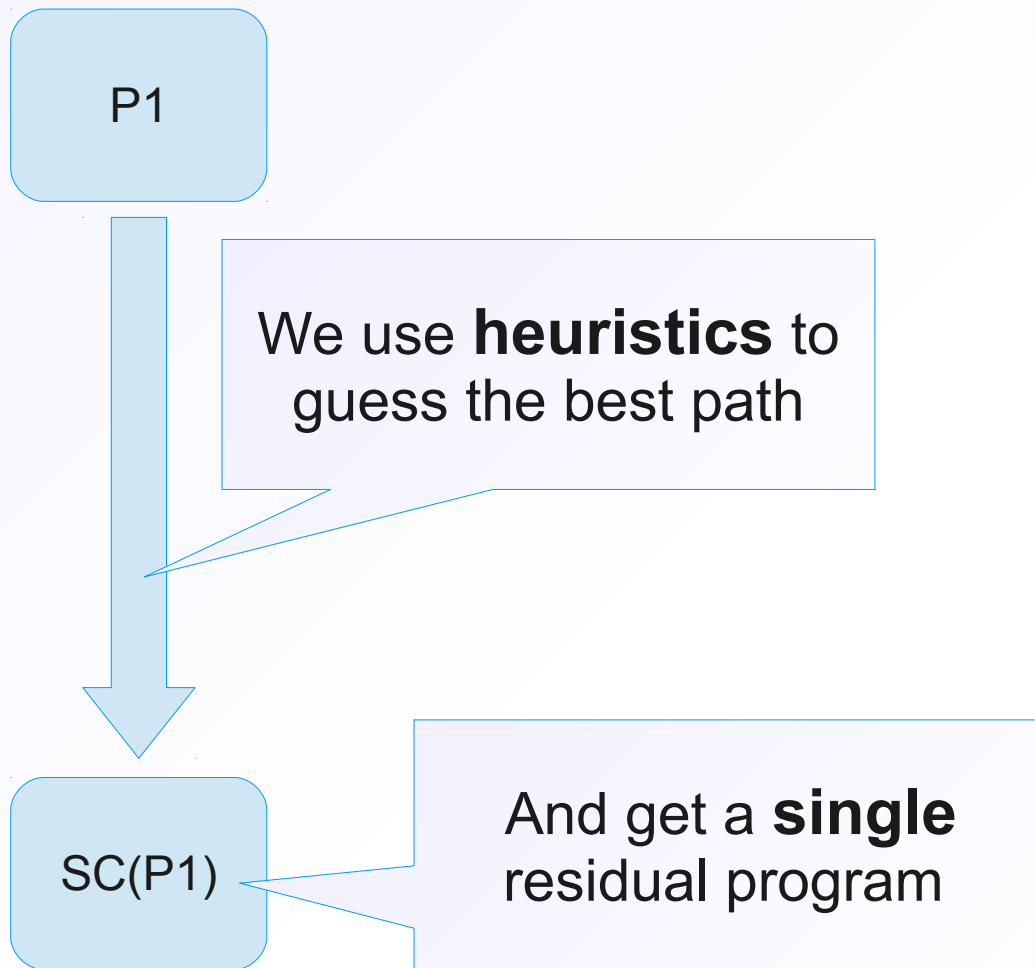
Overgraph Representation for Multi-Result Supercompilation

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Meta 2012

General Idea of Multi-Resulttness



General Idea of Multi-Resulttness



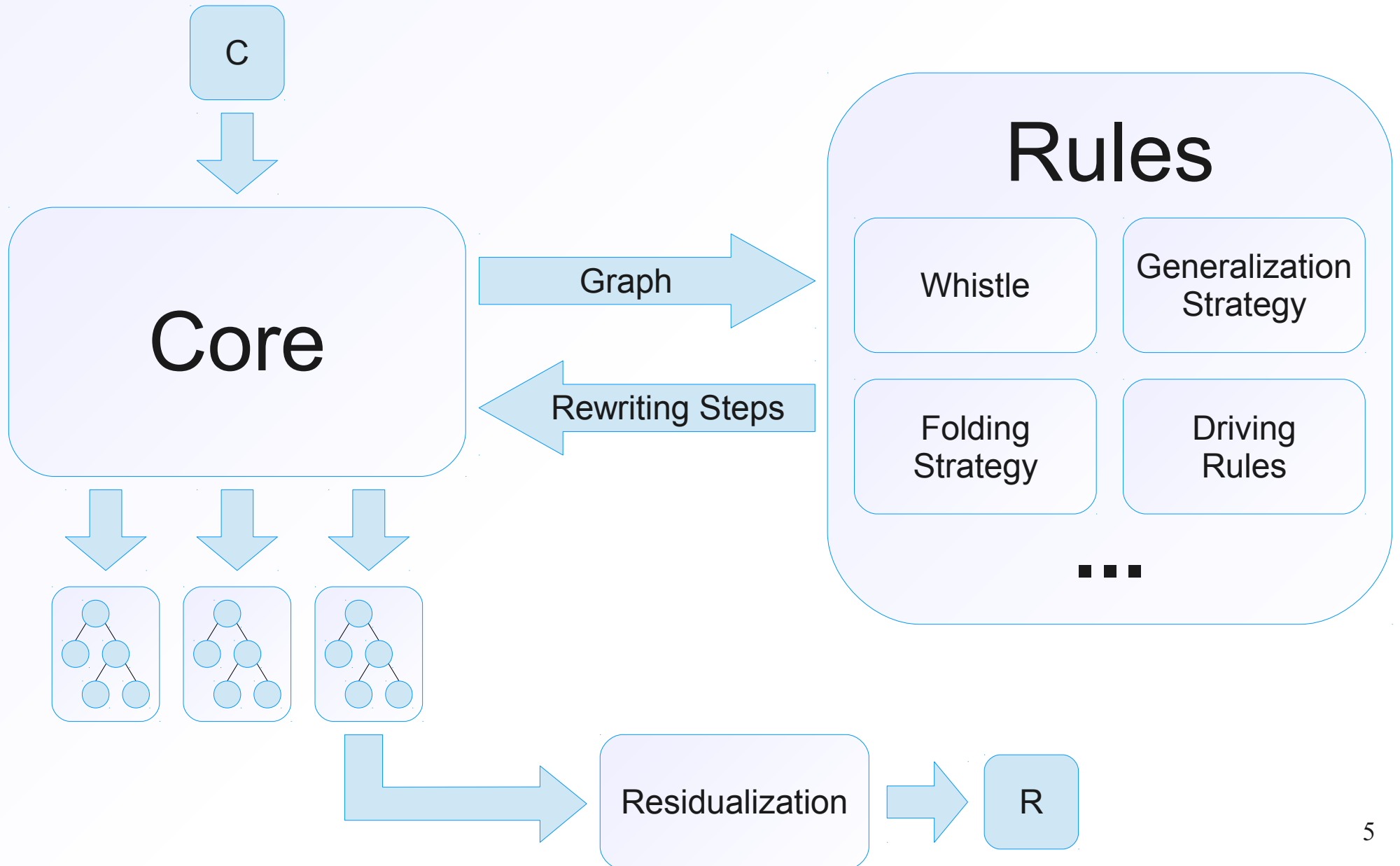
A problem

Millions of residual programs

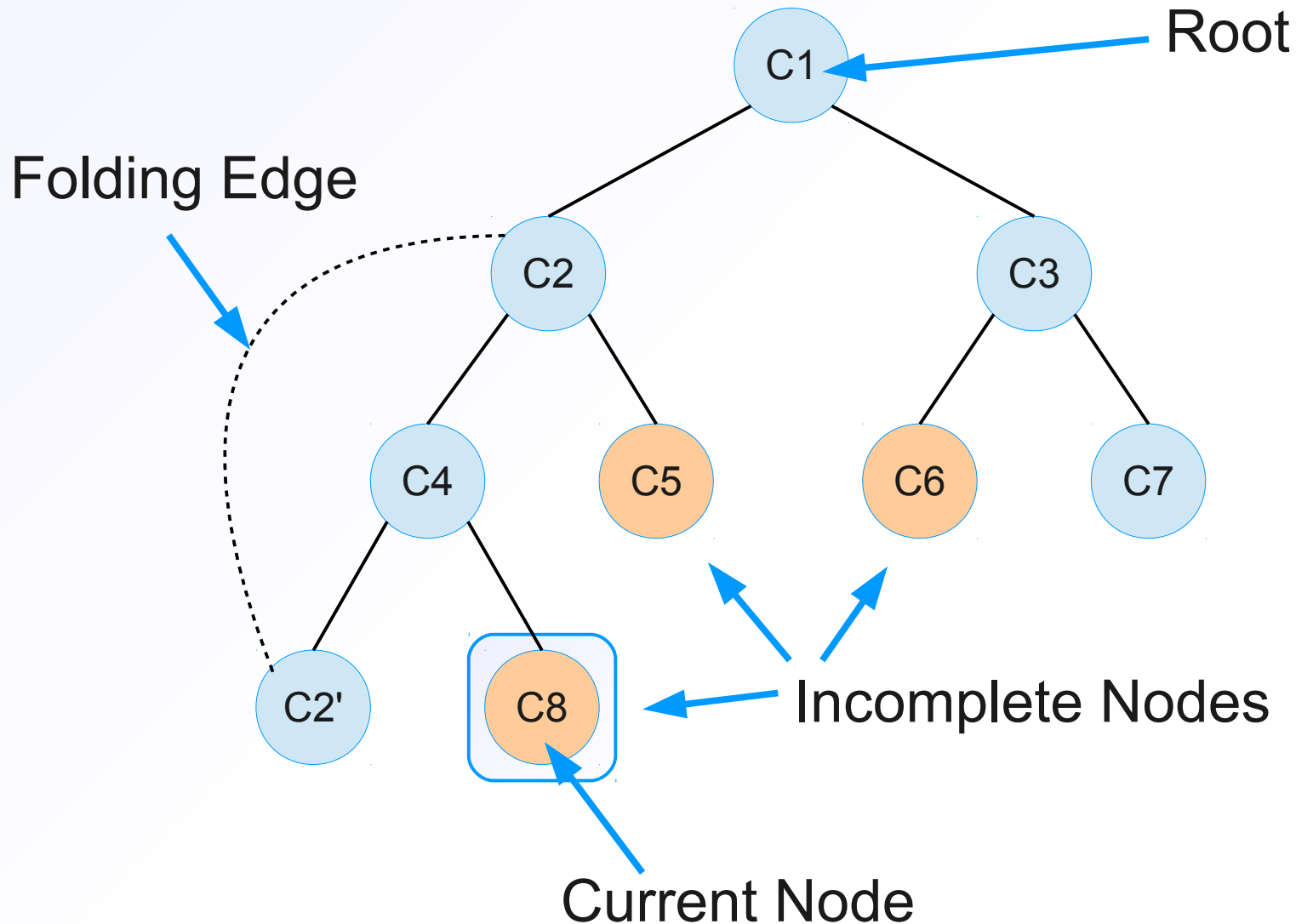
A solution

Overgraph – a compact
representation for sets of graphs

MRSC Toolkit Architecture

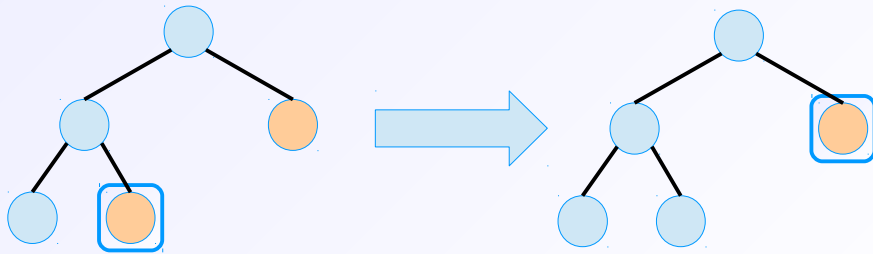


MRSC: Graphs of Configurations

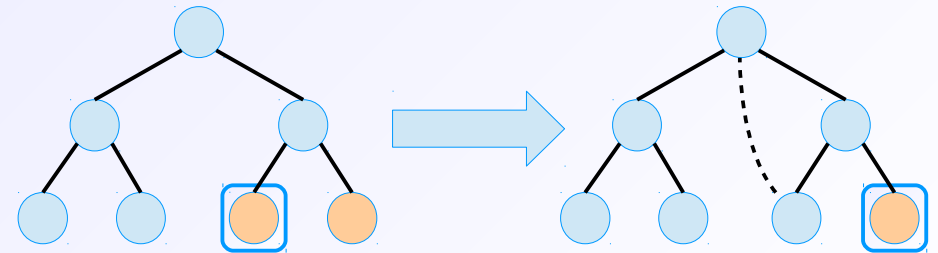


MRSC: Graph Rewriting Steps

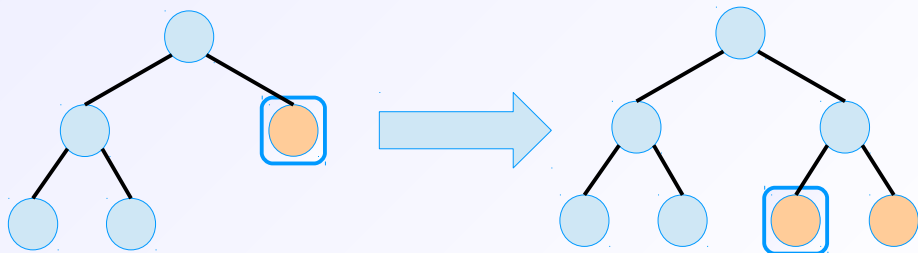
Complete



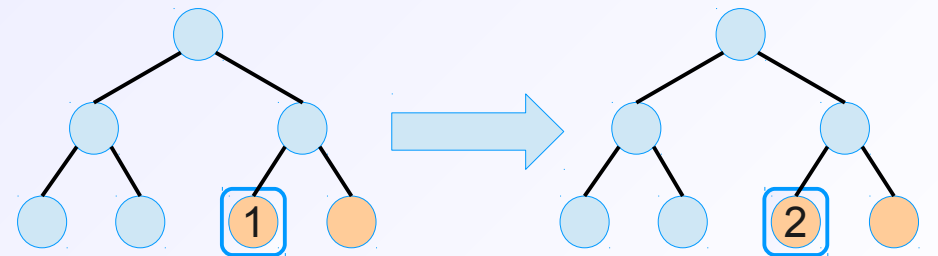
Fold



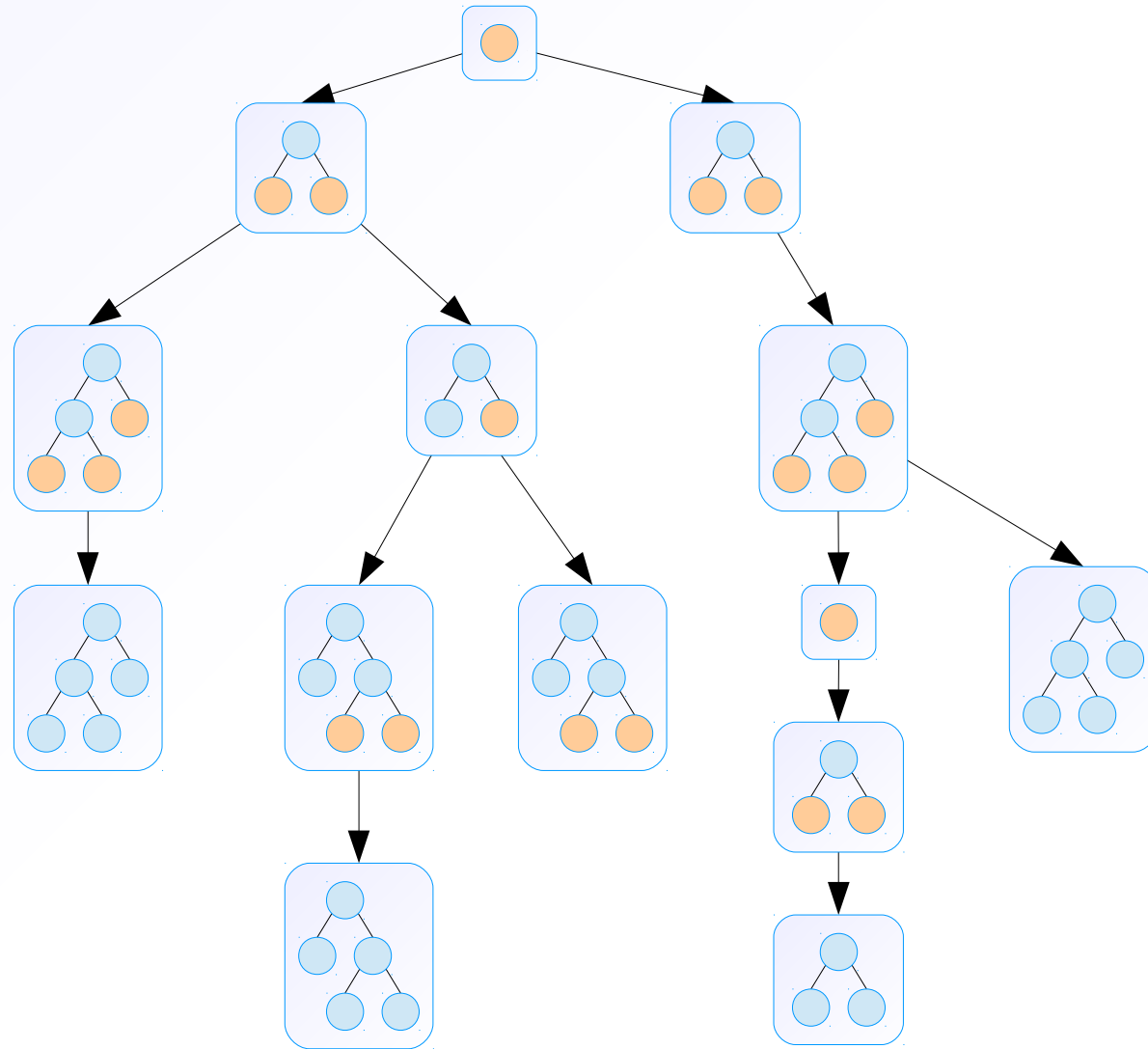
AddChildNodes



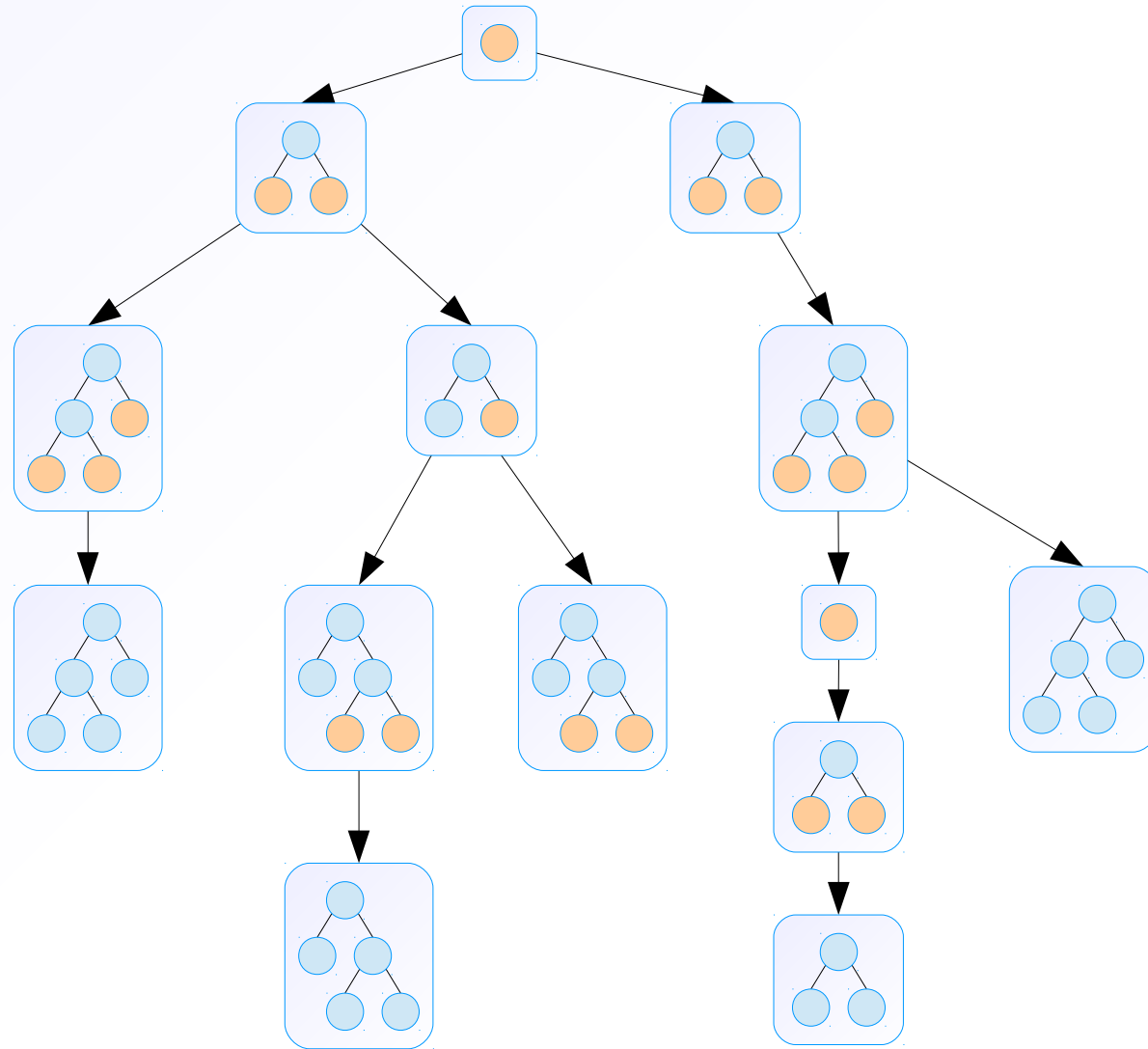
Rebuild



MRSC: Tree of Graphs

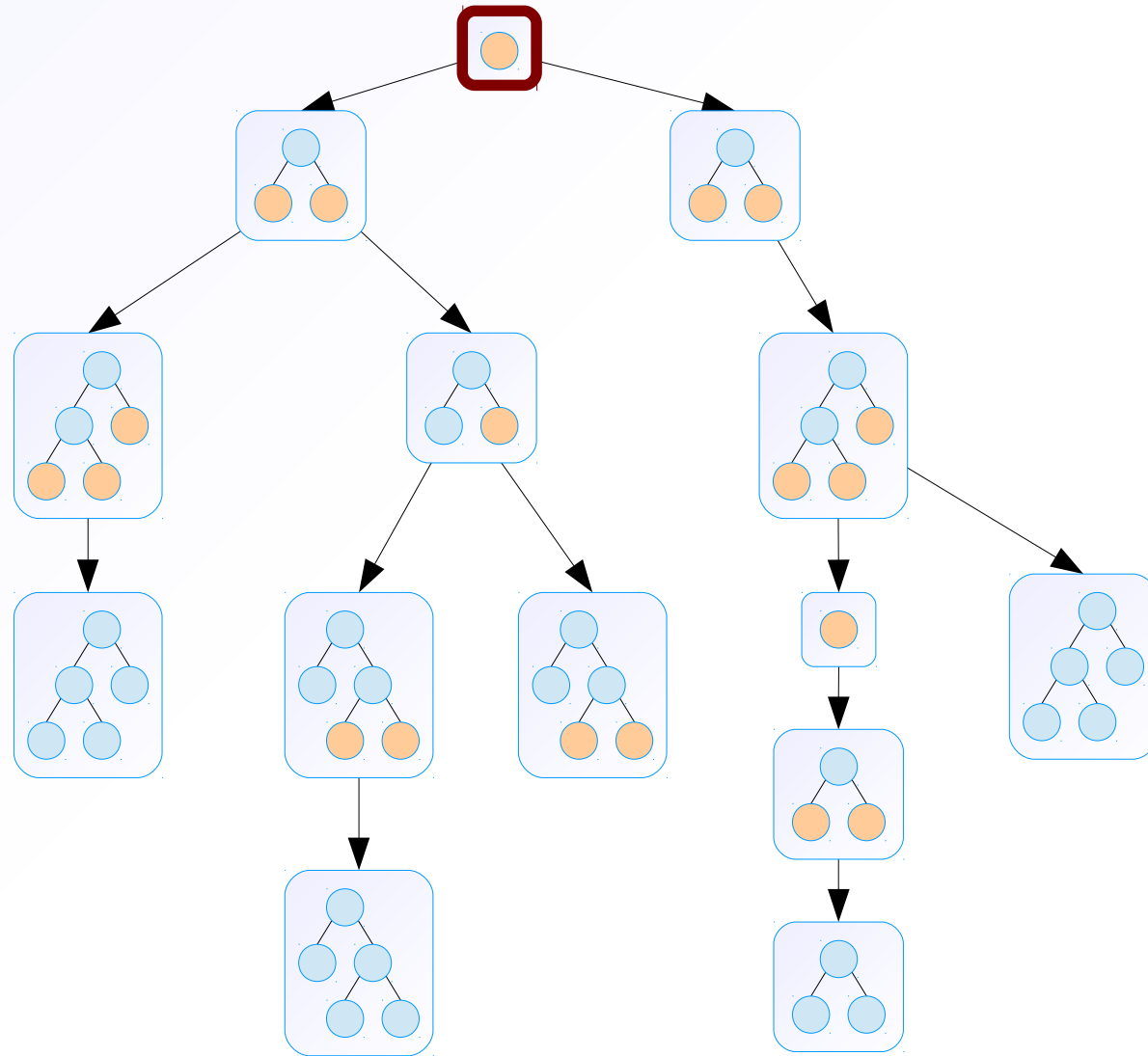


MRSC: Tree of Graphs



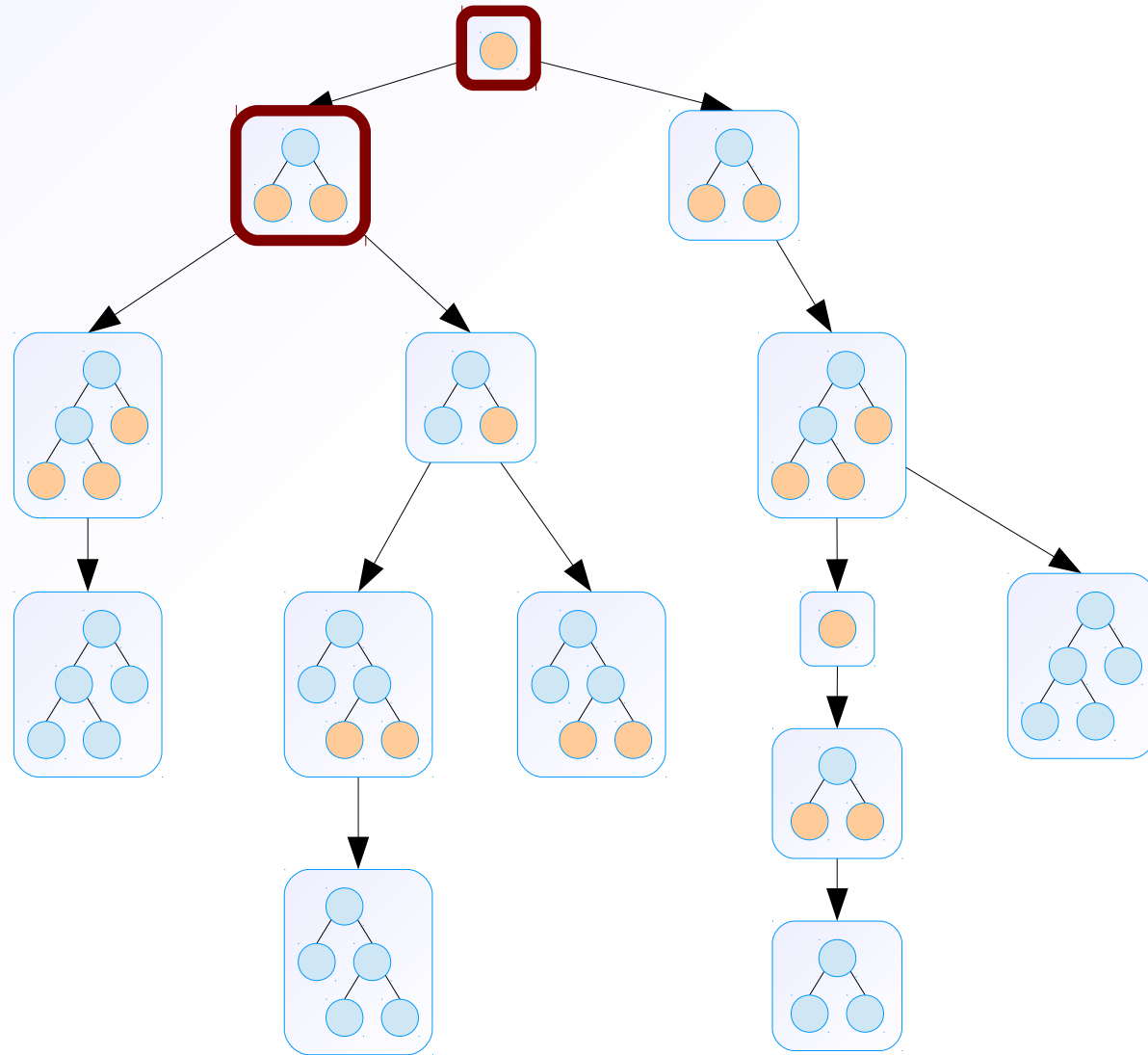
Depth-First Traversal of the Tree of Graphs

MRSC: Tree of Graphs



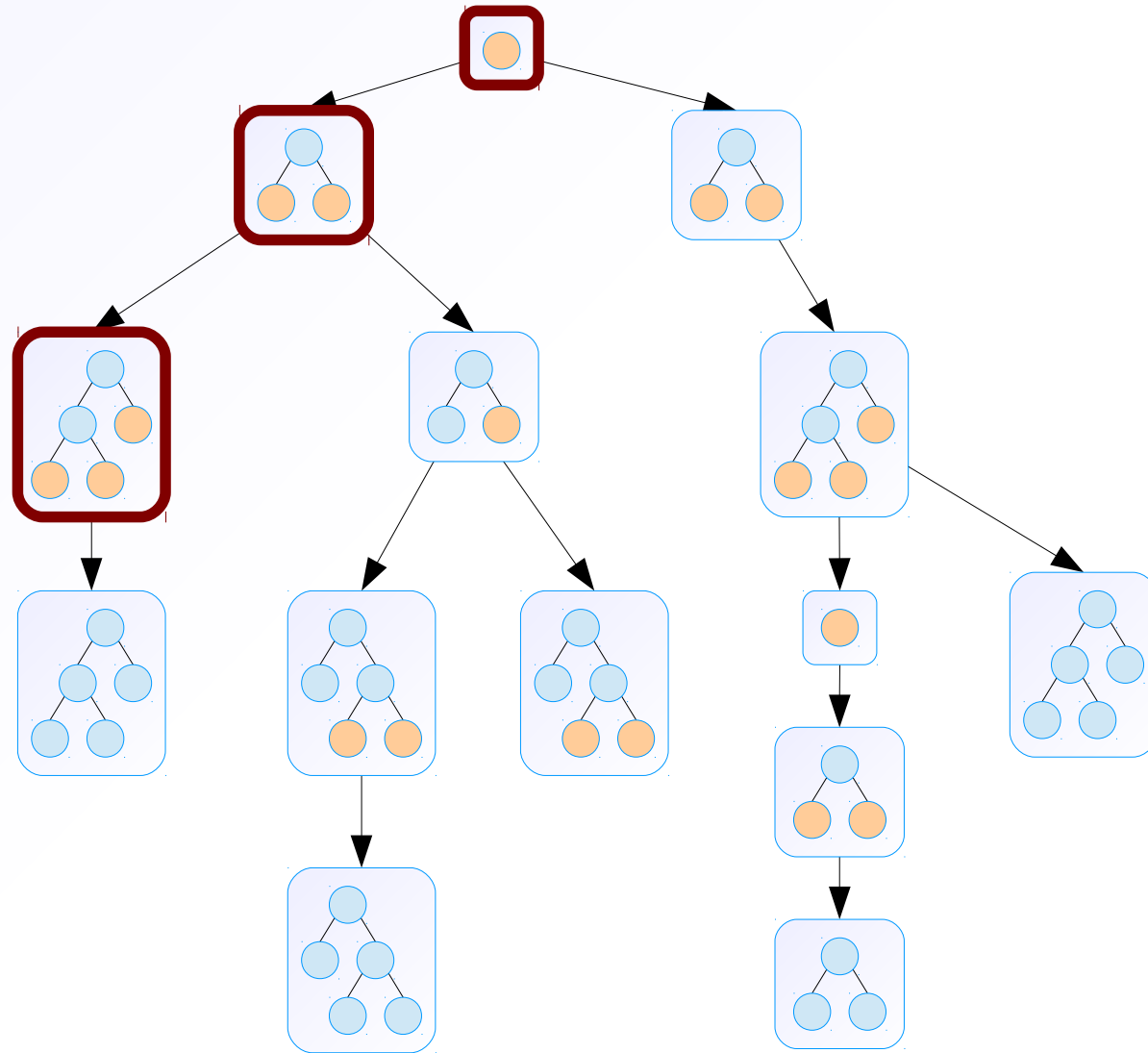
Depth-First Traversal of the Tree of Graphs

MRSC: Tree of Graphs



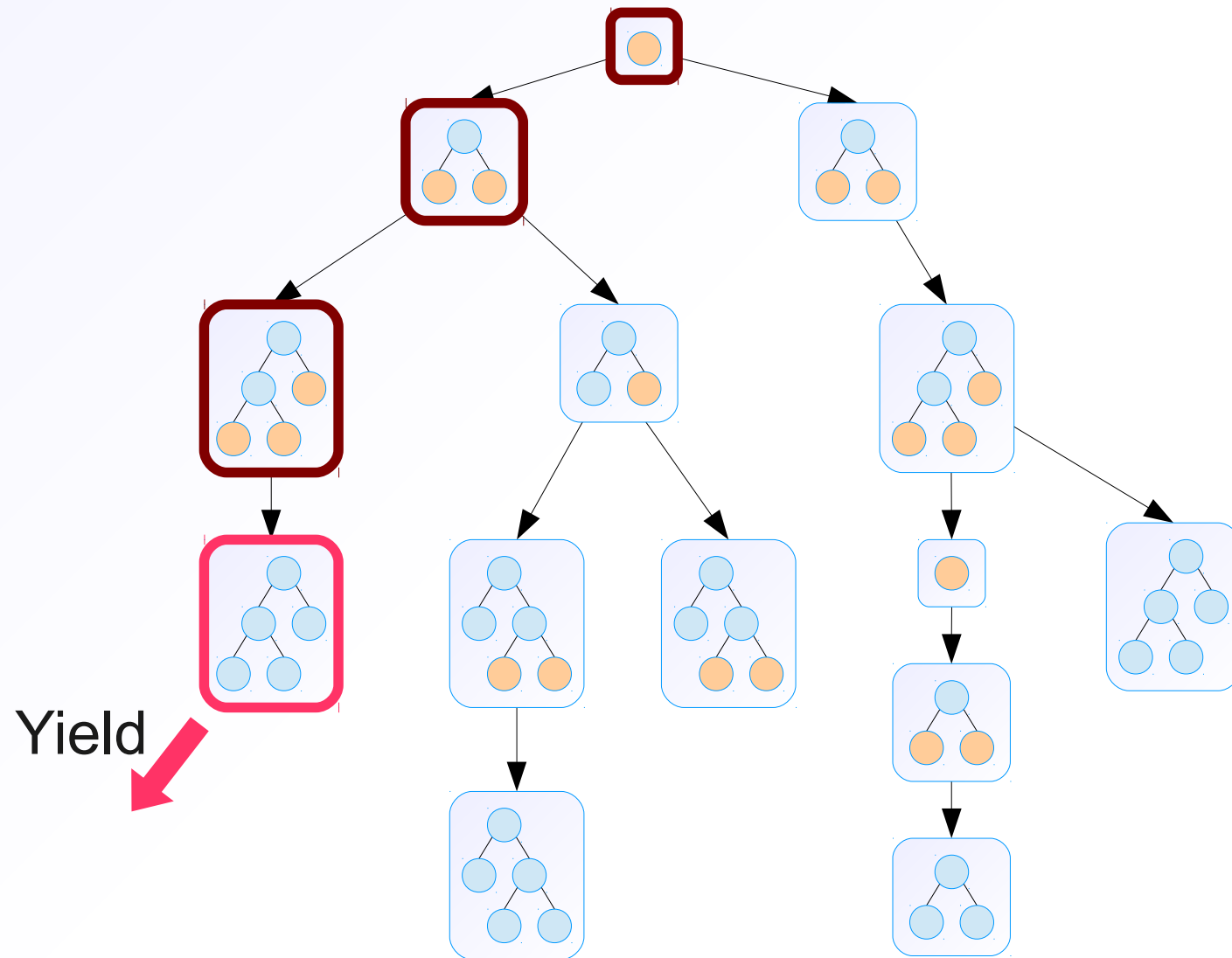
Depth-First Traversal of the Tree of Graphs

MRSC: Tree of Graphs



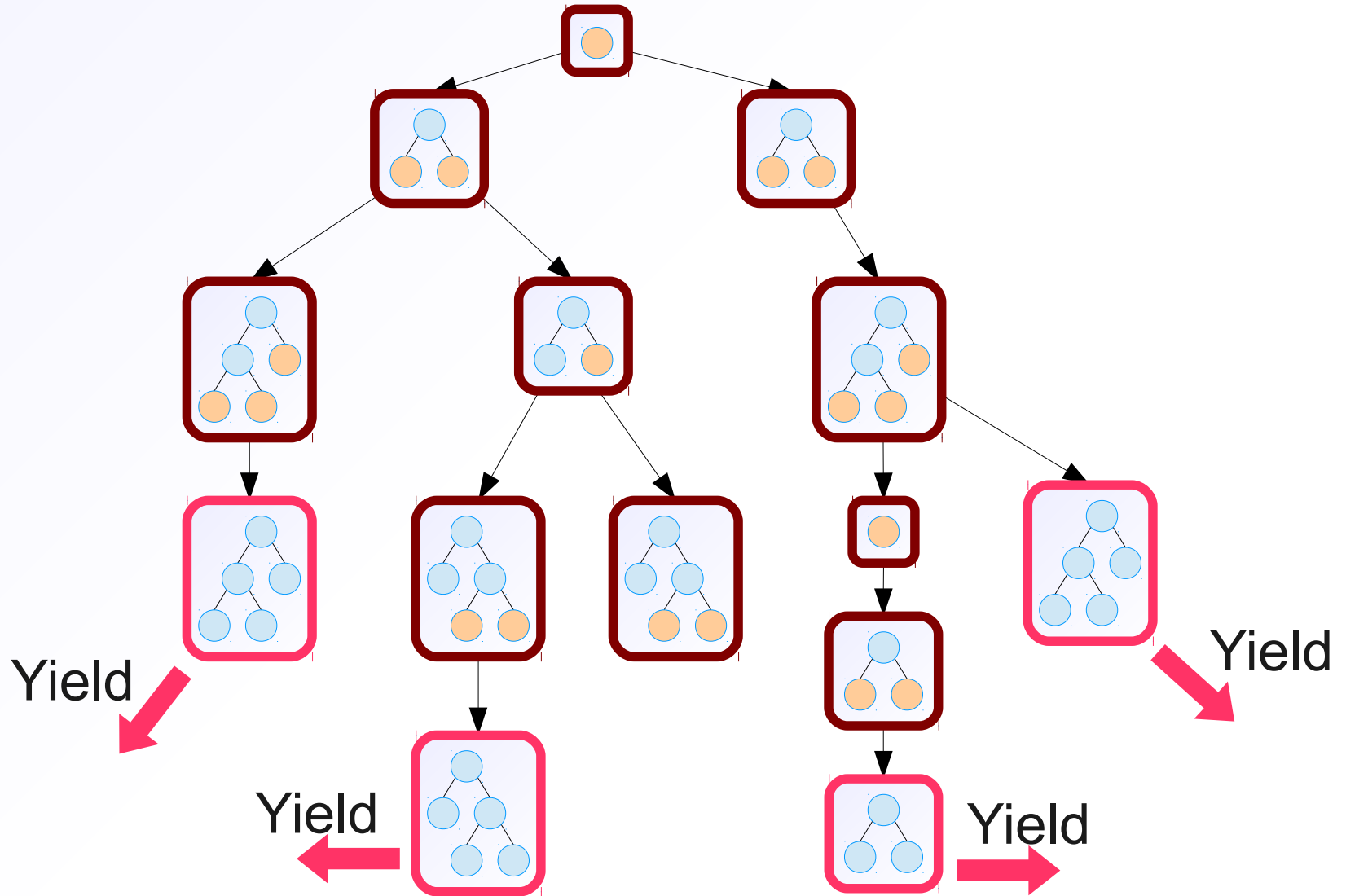
Depth-First Traversal of the Tree of Graphs

MRSC: Tree of Graphs



Depth-First Traversal of the Tree of Graphs

MRSC: Tree of Graphs

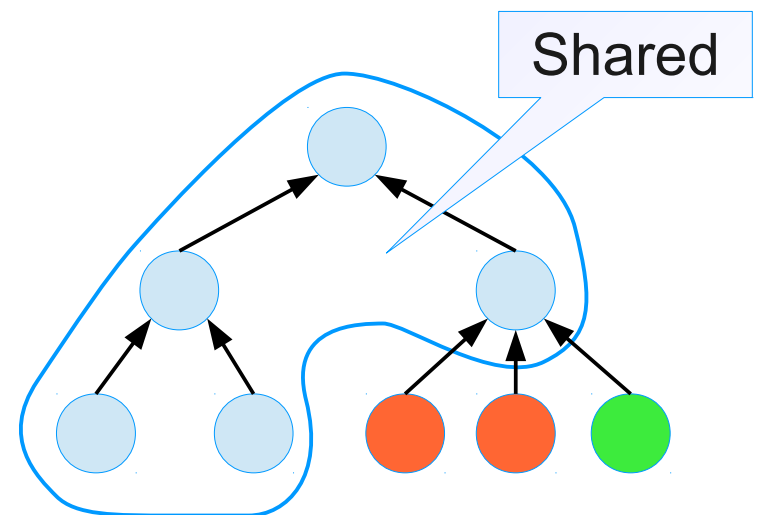
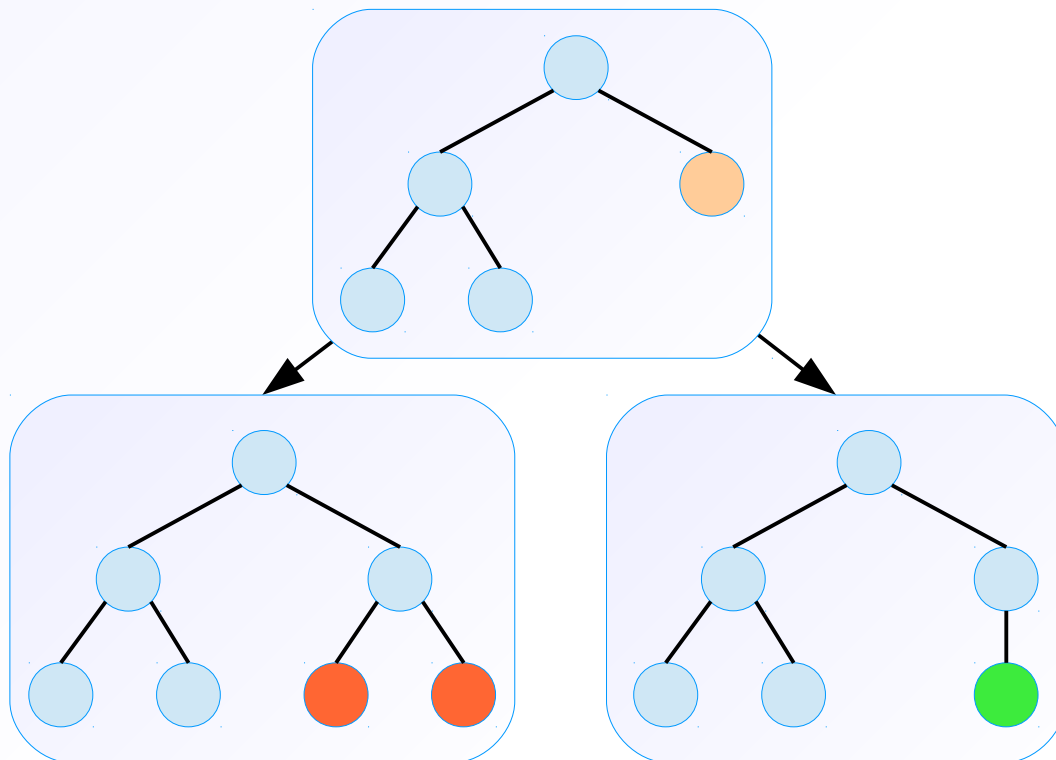


Depth-First Traversal of the Tree of Graphs

Combinatorial Explosion

Too many graphs

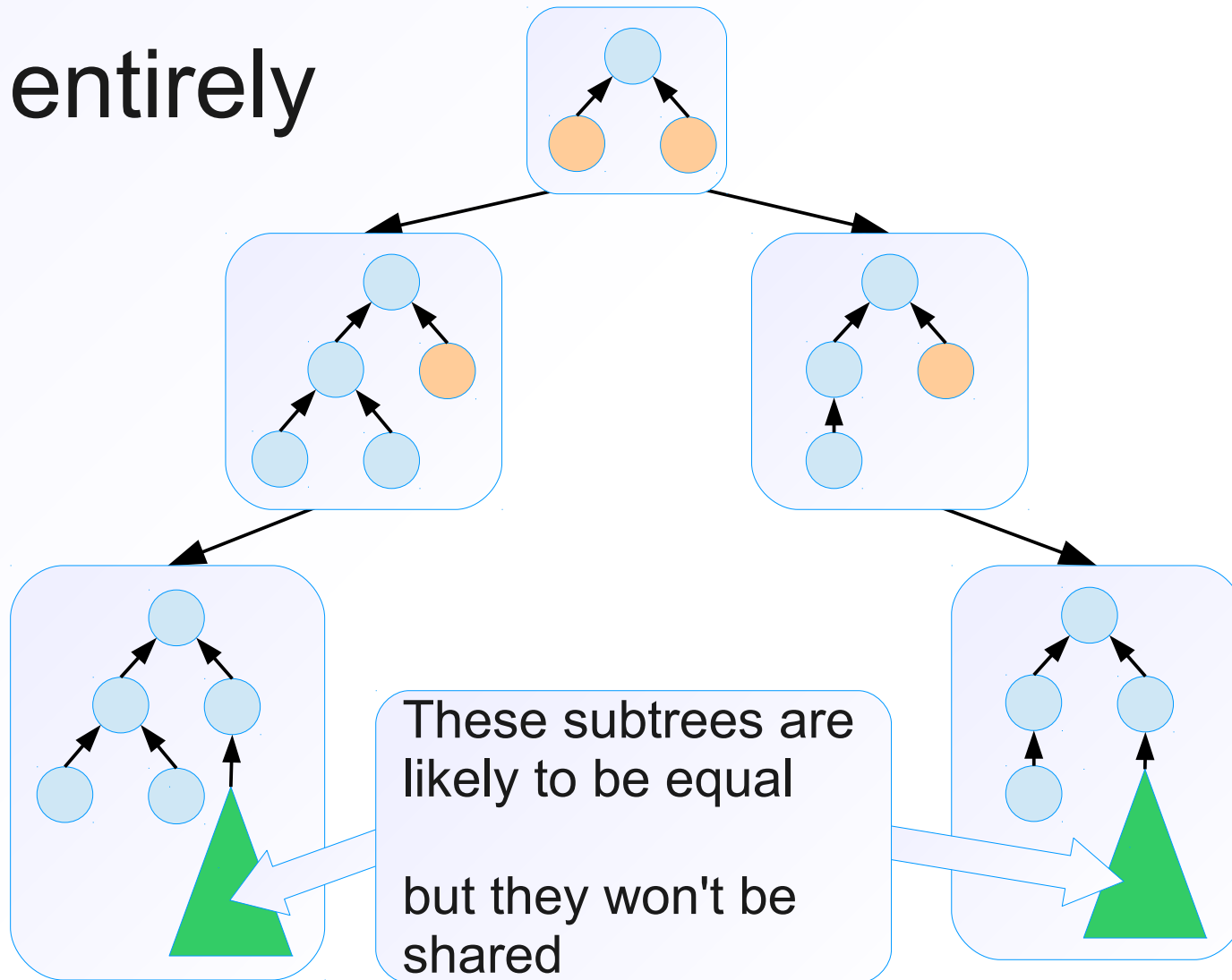
- Use some **heuristics**
- **Share** some parts of graphs



Spaghetti Stack
(MRSC)

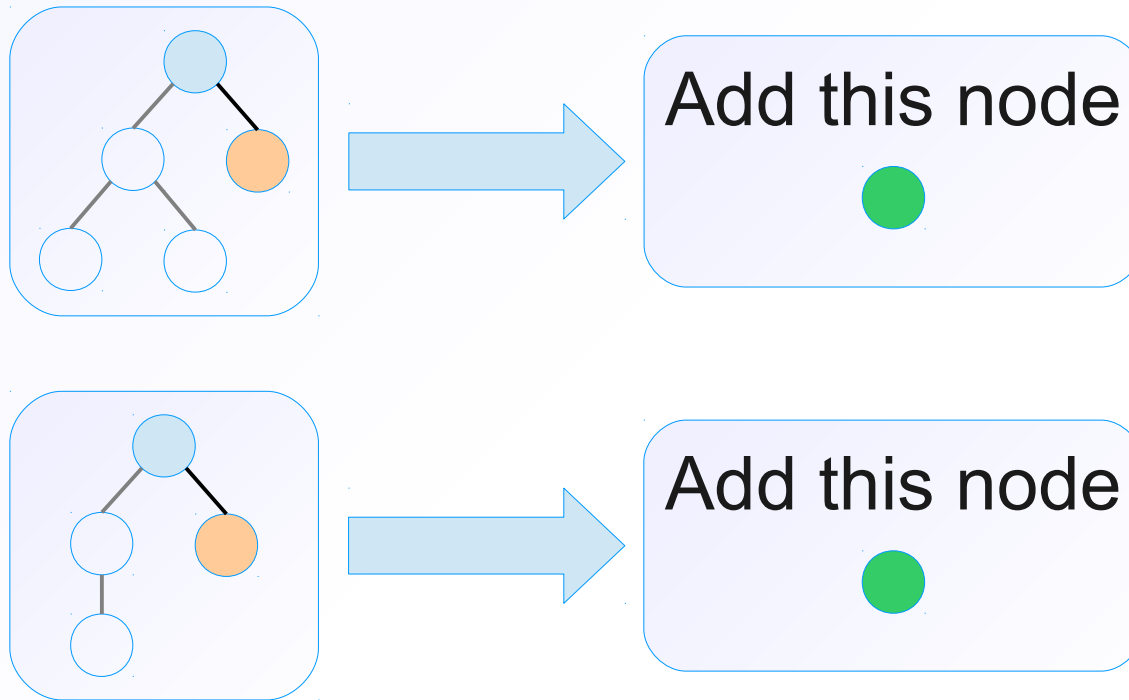
Do Spaghetti Stacks Solve the Problem?

Not entirely



Rules : Graph \rightarrow [Step]

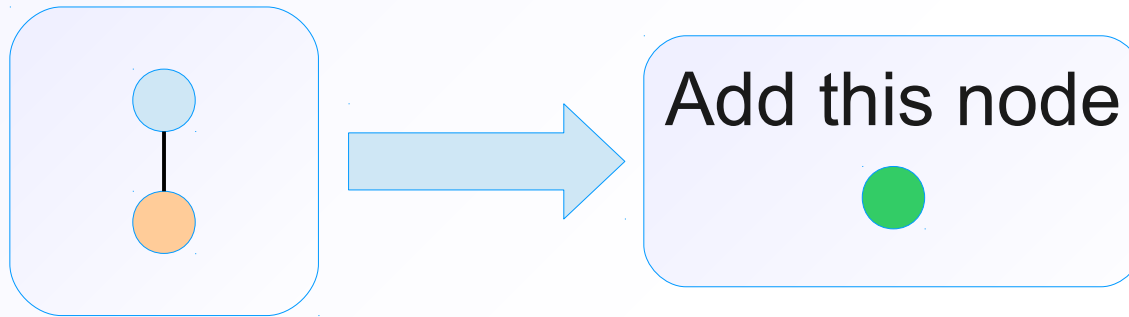
Rules transform graphs into rewriting steps



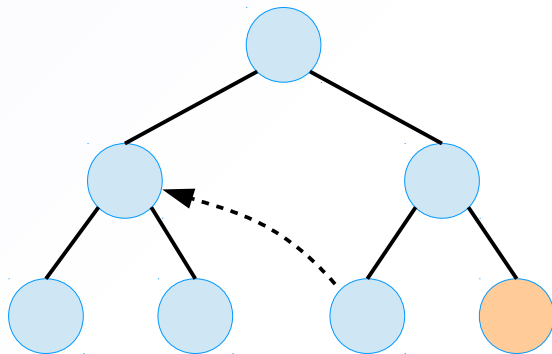
But usually they don't need the whole graph,
just a path from the root to the current node

Rules : Path \rightarrow [Step]

- Let's try to restrict rules to work on paths



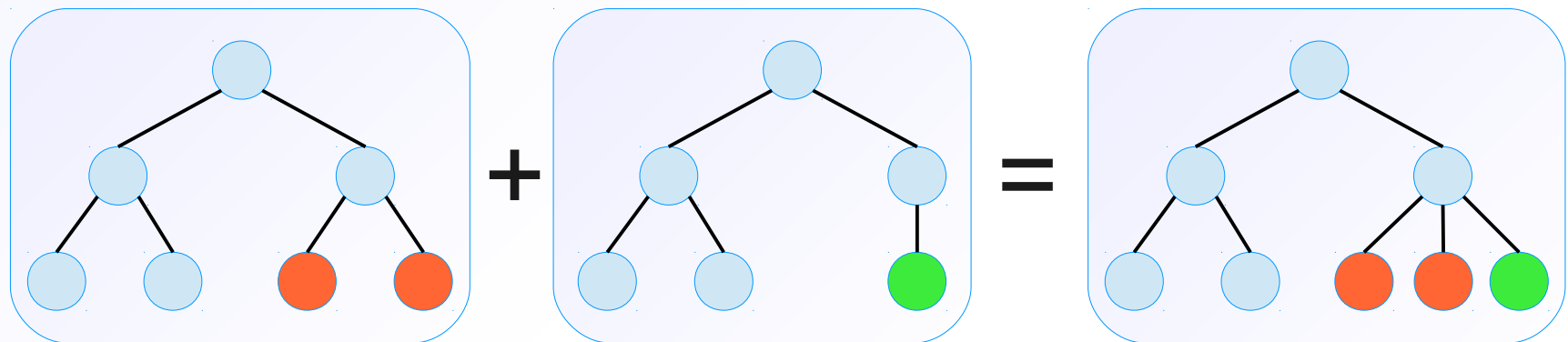
- We would lose an interesting ability to fold with cross edges



- We would need some new **representation** to make use of this new property

Overtree Representation

Let's combine all configuration trees into one big overtree



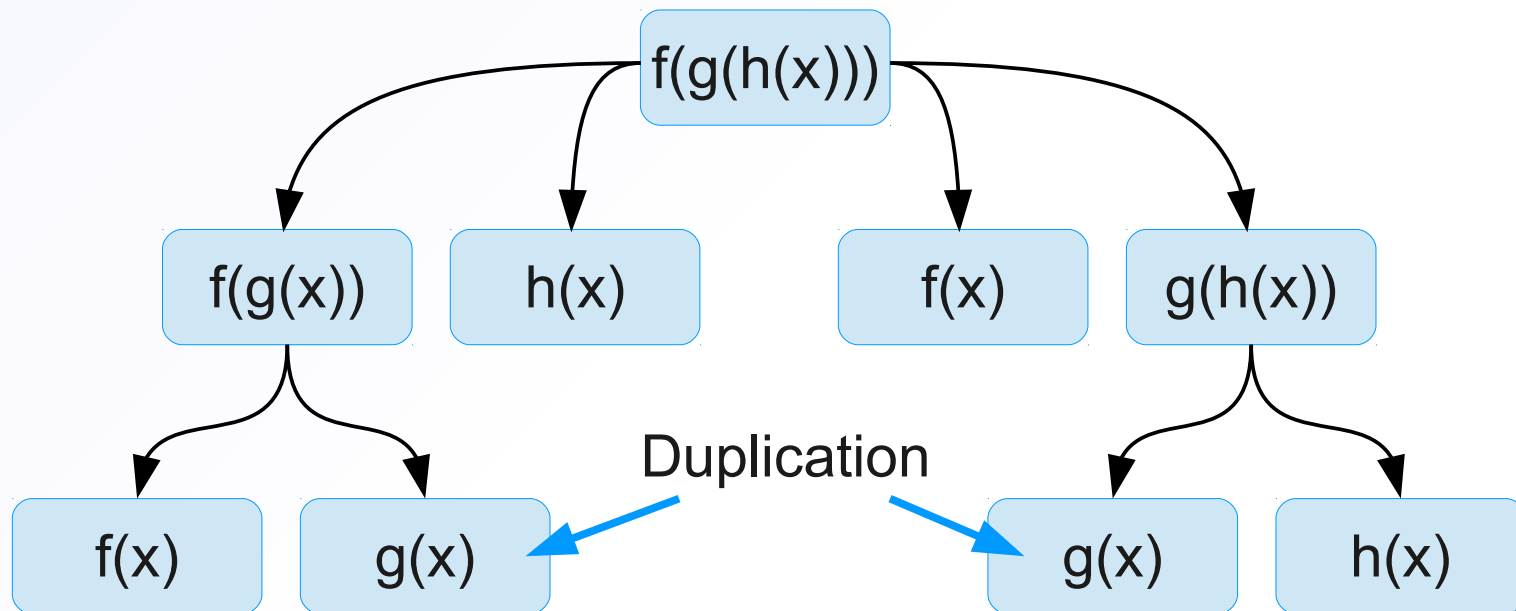
An overtree represents a **set** of trees

```
data Tree = Tree (F Tree)
```

```
data OTree = OTree [F OTree]
```

Do Overtrees Solve the Problem?

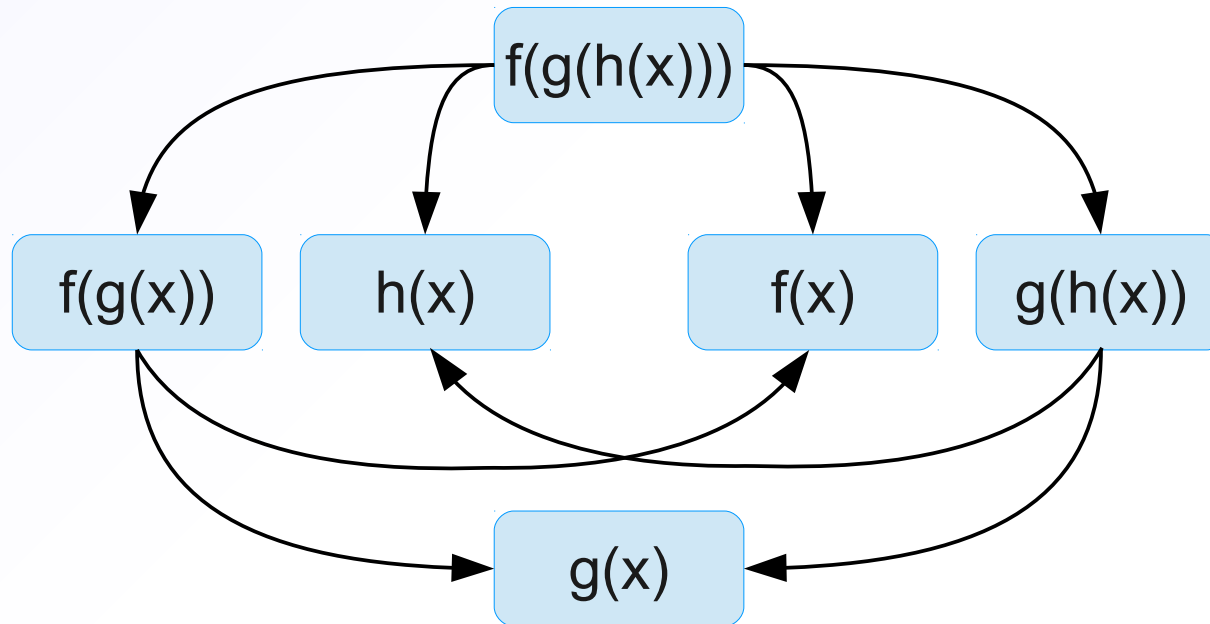
- They are a bit better, but still...



- We've already lost cross edges
- Are we going to lose folding edges completely?

Overgraph

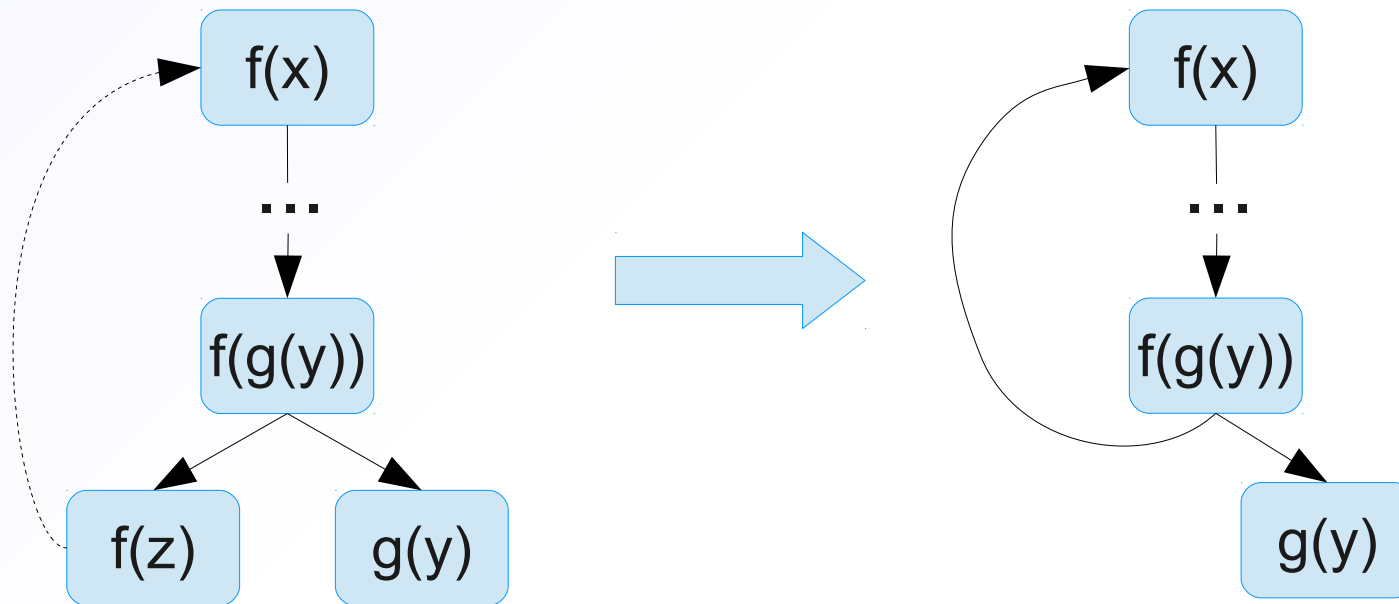
- Let's just glue together nodes equivalent **up to renaming**



- Each configuration corresponds to no more than one node

Folding

We don't need special folding edges

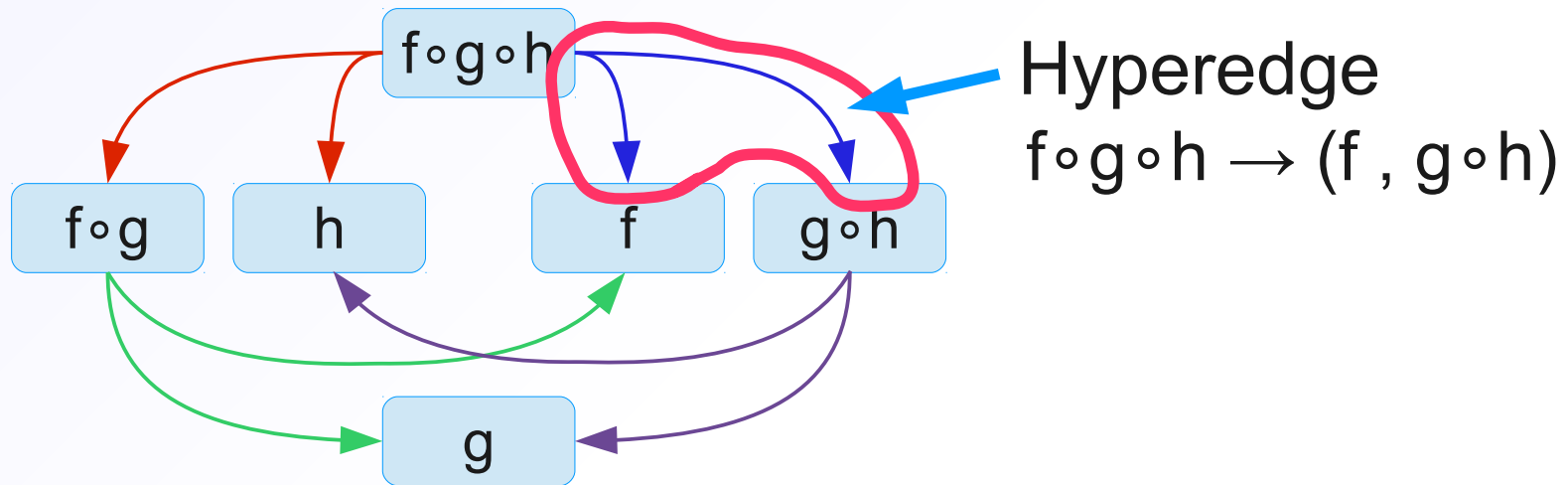


Advantages and Problems

- Overgraphs are more compact
- Overgraphs are cleaner
 - One configuration — one node
 - No special folding edges
- Overgraphs contain more information
- Each node can have multiple parents
 - Can we use binary whistles?
 - How can we control generalization?
- How to apply rules?
- How to extract residual programs?

Hyperedges

- We will call **bundles of edges** hyperedges



- Hyperedges represent steps like driving and generalization
- Completion step can be represented as a hyperedge with **zero destination nodes**

C1

$C1 \rightarrow ()$

C2

incomplete nodes have no outgoing hyperedges

Supercompilation with Overgraphs

1) Overgraph **Construction**

Add nodes and edges while possible

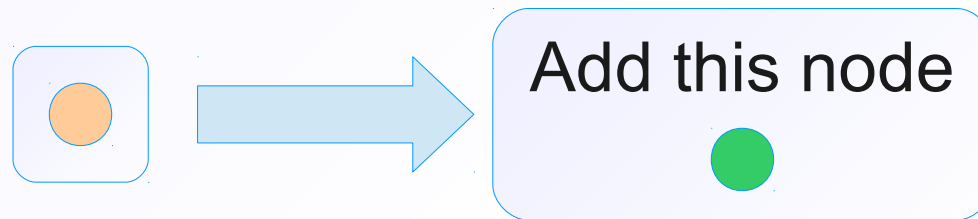
2) Overgraph **Truncation**

Remove useless nodes and edges

3) **Residualization**

Overgraph Construction

- Rule : Configuration \rightarrow [Step]



- Rule : Overgraph \rightarrow [Hyperedge]

In what order should we apply the rules?

r is monotone if for all graphs G and H :

$$G \subseteq H \Rightarrow r(G) \subseteq r(H)$$

If all rules are monotone we can apply them in any order

Rules

- We can also write rules in this form:

$$\frac{\text{precondition}}{\text{hyperedges to add}}$$

- Examples:

$$\frac{\neg \text{UnaryWhistle}(c)}{c \rightarrow \text{drive}(c)}$$

$$\frac{\text{always}}{c \rightarrow \text{generalize}(c)}$$

$$\frac{\text{min_depth}(c) < 42}{c \rightarrow \text{drive}(c)}$$

This precondition is monotone

Binary Whistles

$\neg \exists d \in G : \text{BinaryWhistle}(c,d)$

NOT monotone

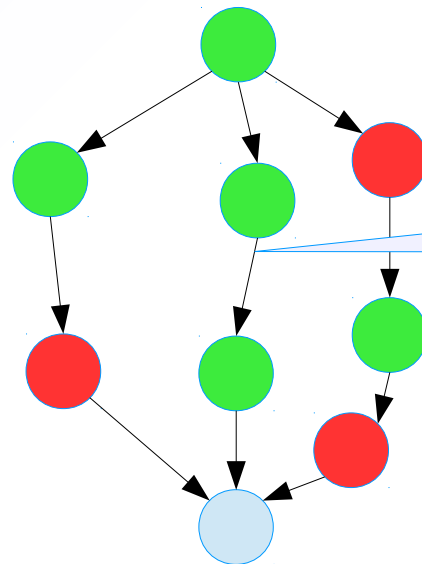
$c \rightarrow \text{drive}(c)$

\exists path p from root to c :

$\forall d \in p : \neg \text{BinaryWhistle}(c,d)$

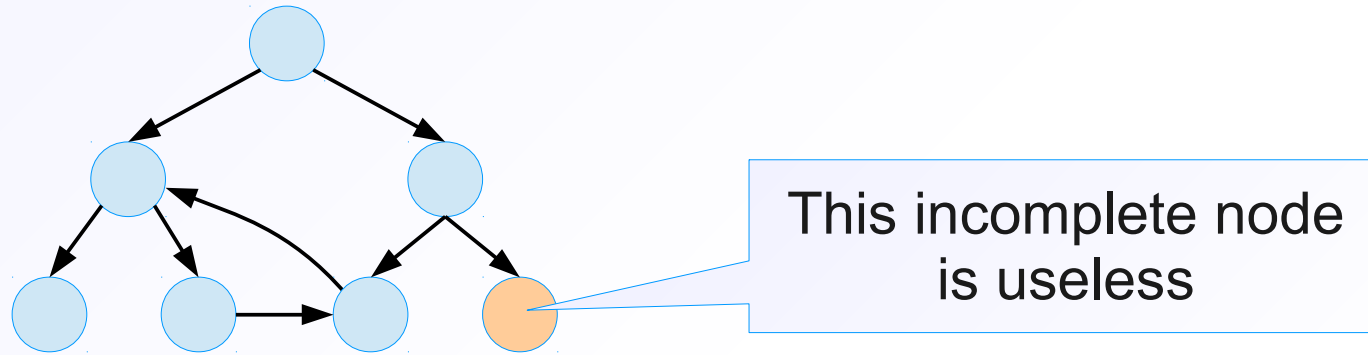
OK

$c \rightarrow \text{drive}(c)$

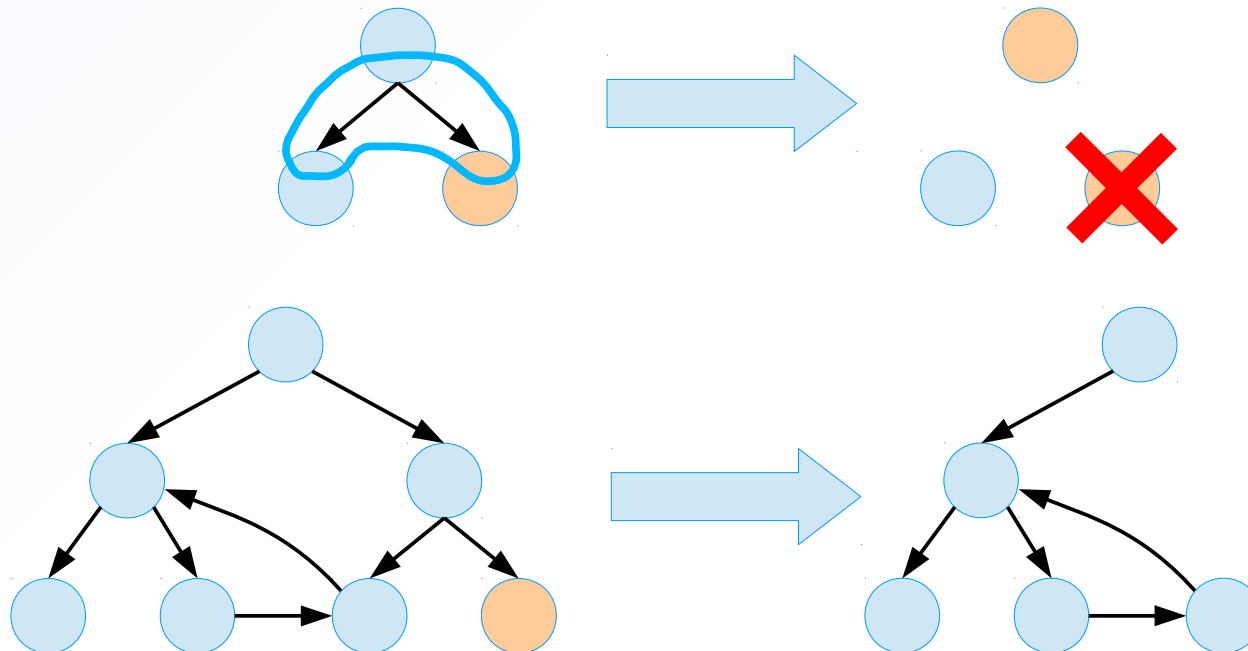


This green path won't disappear

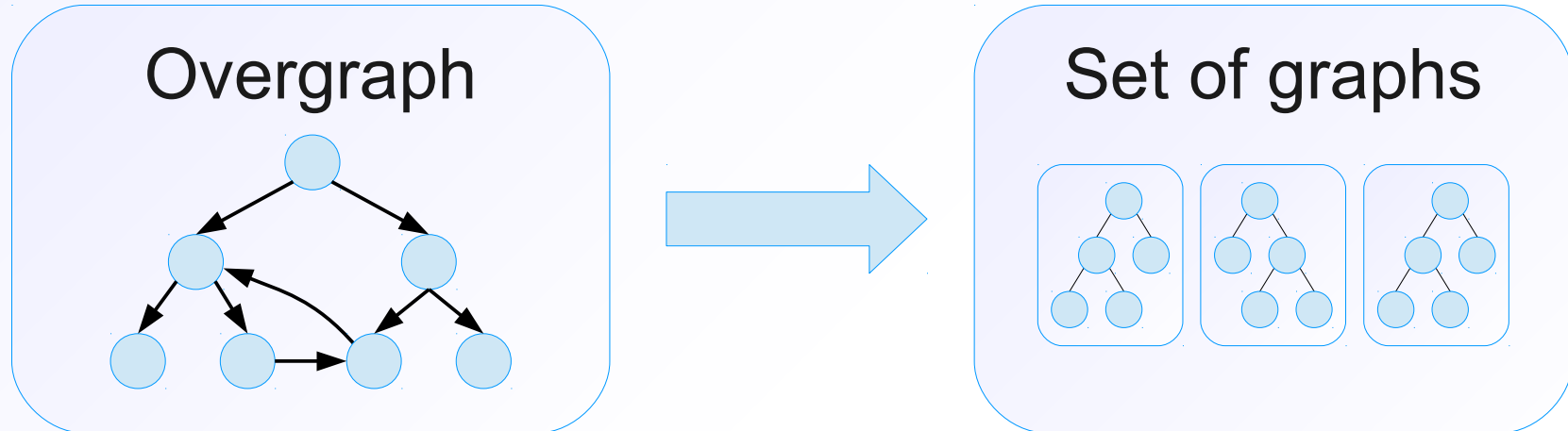
Overgraph Truncation



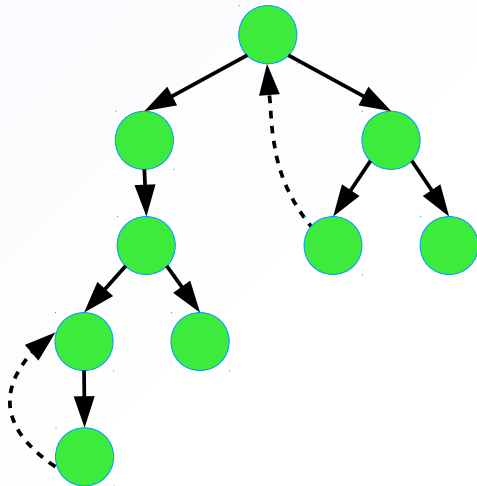
We should remove all incident hyperedges



Residualization

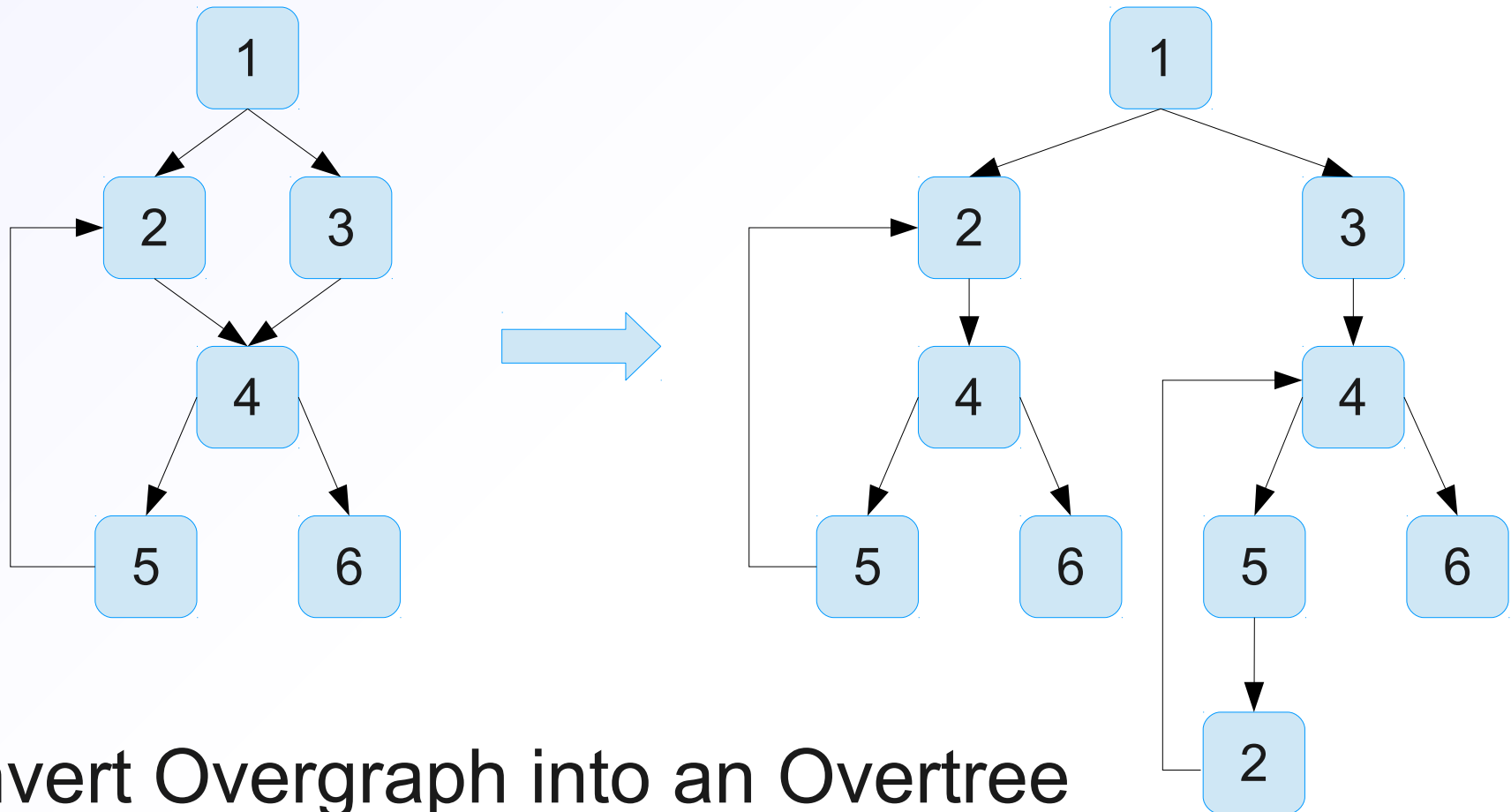


Building a full set of graphs should be avoided!



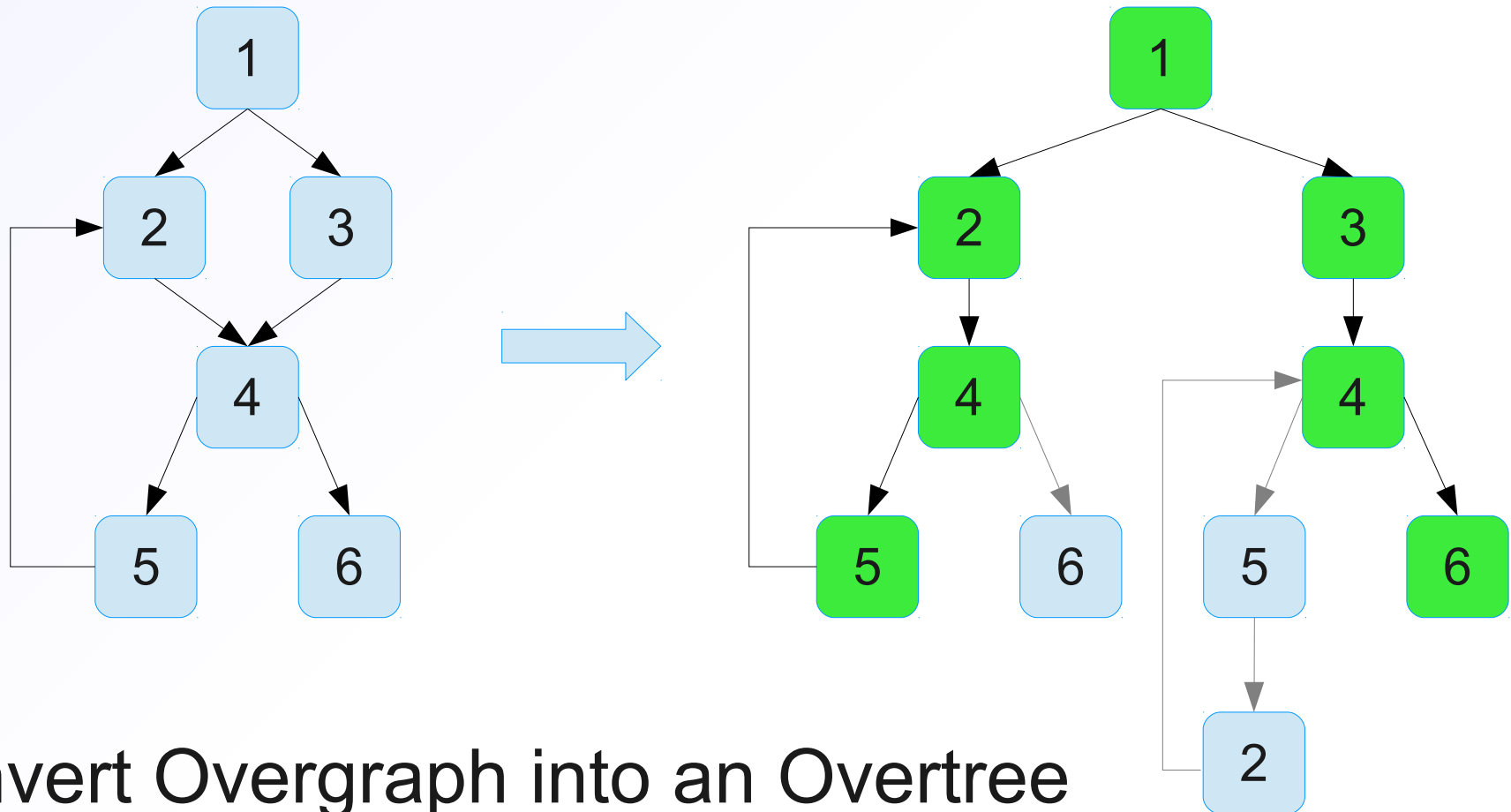
We will represent residual programs as **trees with back edges** (i.e. no subprogram sharing)

Naive Residualization Algorithm



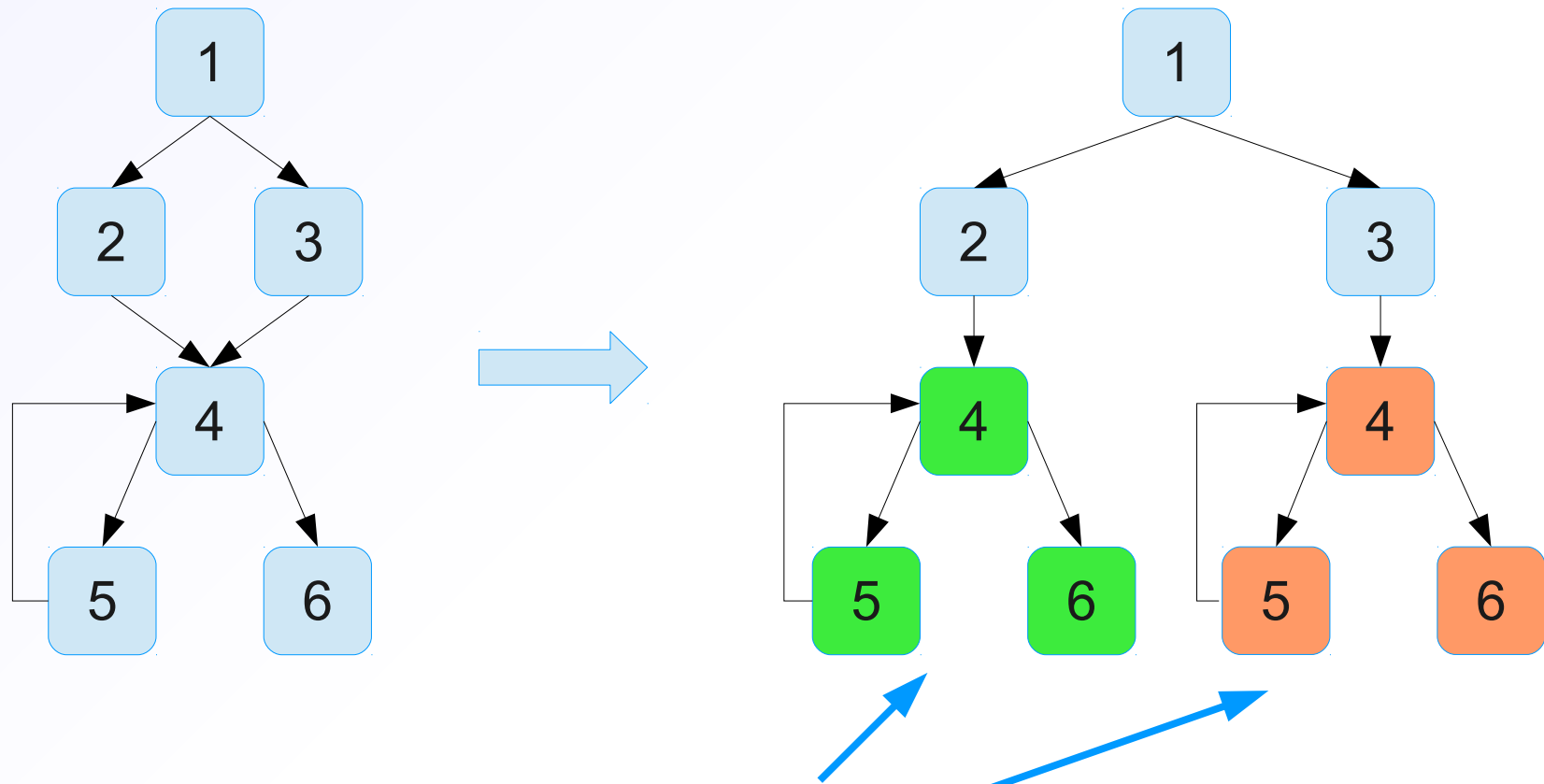
Convert Overgraph into an Overtree
and then convert it into a set of trees

Naive Residualization Algorithm



Convert Overgraph into an Overtree
and then convert it into a set of trees

Suboptimality



Absolutely identical subtrees

Idea: Cache intermediate results

More Formal Definition

$R : \text{Node} \rightarrow [\text{Node}] \rightarrow [\text{Tree}]$

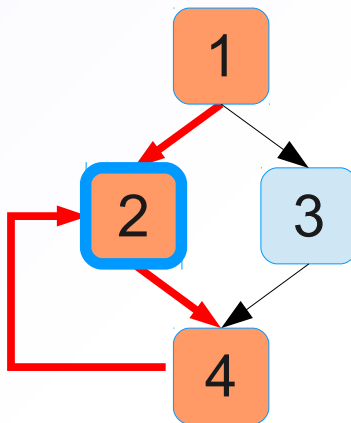
$R \ n \ h \mid n \in h = [\text{Fold}(n)]$

$R \ n \ h \mid \text{otherwise} =$

$[n \rightarrow (r1 \dots rk) \mid$

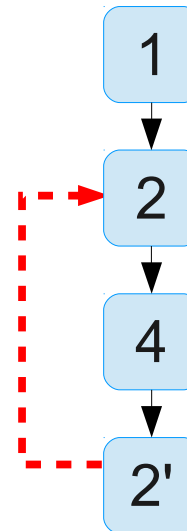
$n \rightarrow (d1 \dots dk) \in G,$

$ri \in R \ di \ (n:h)]$



$h = [4, 2, 1]$

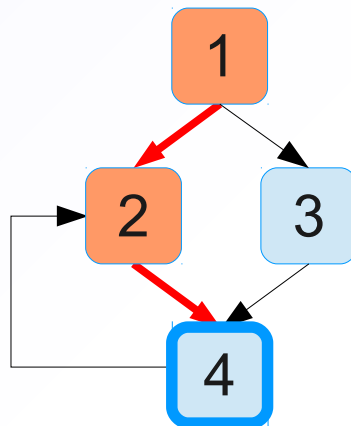
$n = 2$



More Formal Definition

$R : \text{Node} \rightarrow [\text{Node}] \rightarrow [\text{Tree}]$

```
R n h | n ∈ h = [Fold(n)]  
R n h | otherwise =  
  [n → (r1 ... rk) |  
    n → (d1 ... dk) ∈ G,  
    ri ∈ R di (n:h)]
```



$h = [2, 1]$

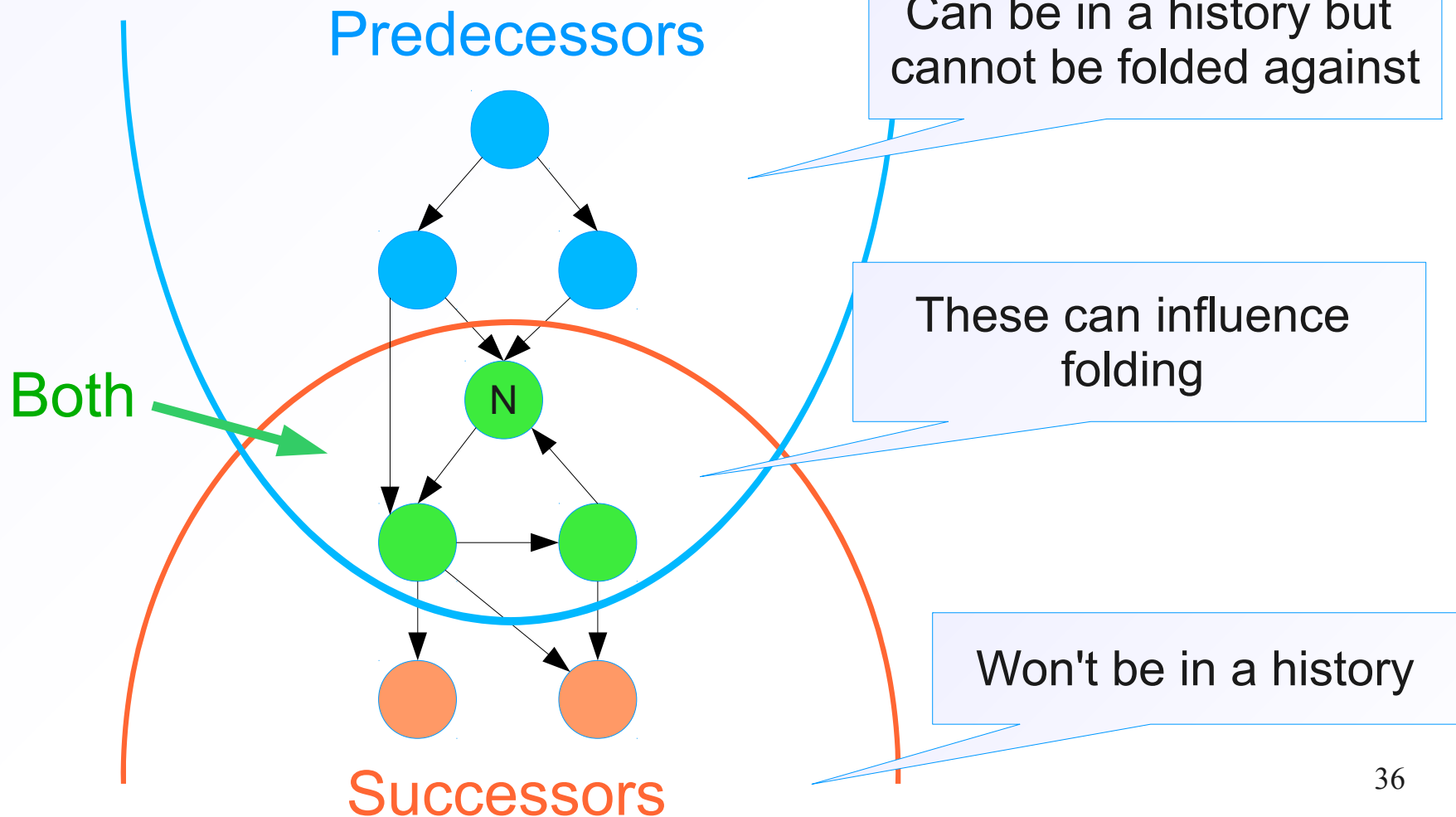
$n = 4$

$R\ 2\ [4, 2, 1]$

History Structure

R : Node \rightarrow [Node] \rightarrow [Tree]

History \nearrow



Enhanced Residualization

- Removing pure predecessors from history won't change the result

$$R \ n \ h = R \ n \ (h \ n \ \text{succs}(n))$$

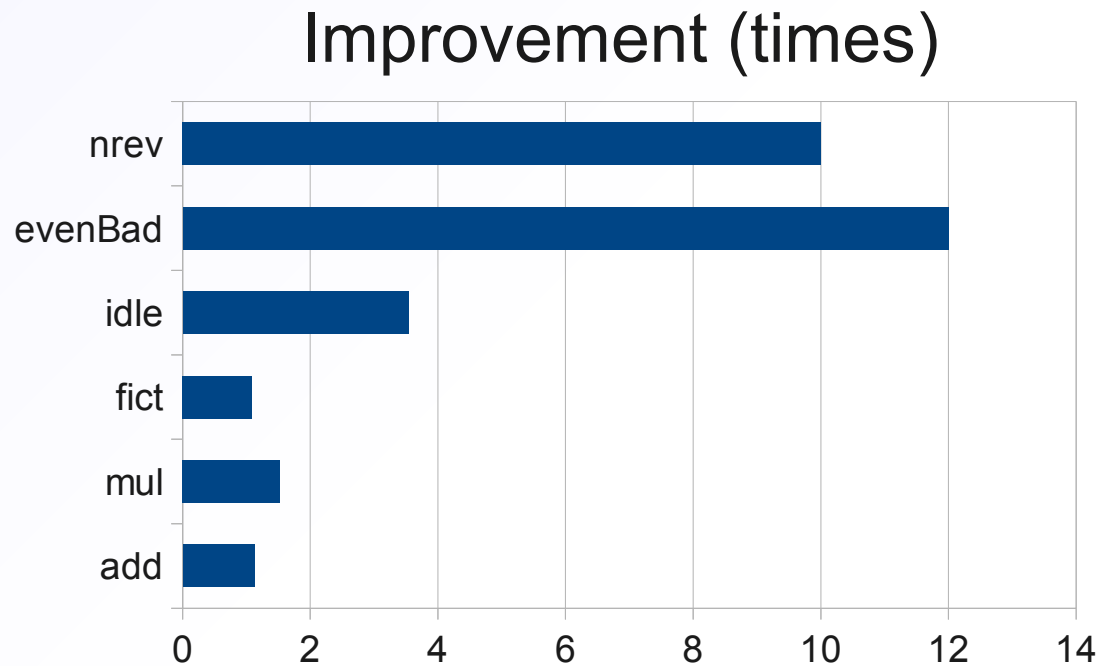
- Let's rewrite residualization algorithm this way:

```
R n h | n ∈ h = [Fold(n)]
R n h | otherwise =
  [n → (r1 ... rk) |
    n → (d1 ... dk) ∈ G,
    ri ∈ R di (n:h n succs(di))]
```

- Now we can just apply memoization

Evaluation of Residualization Algorithms

- Caching improves performance

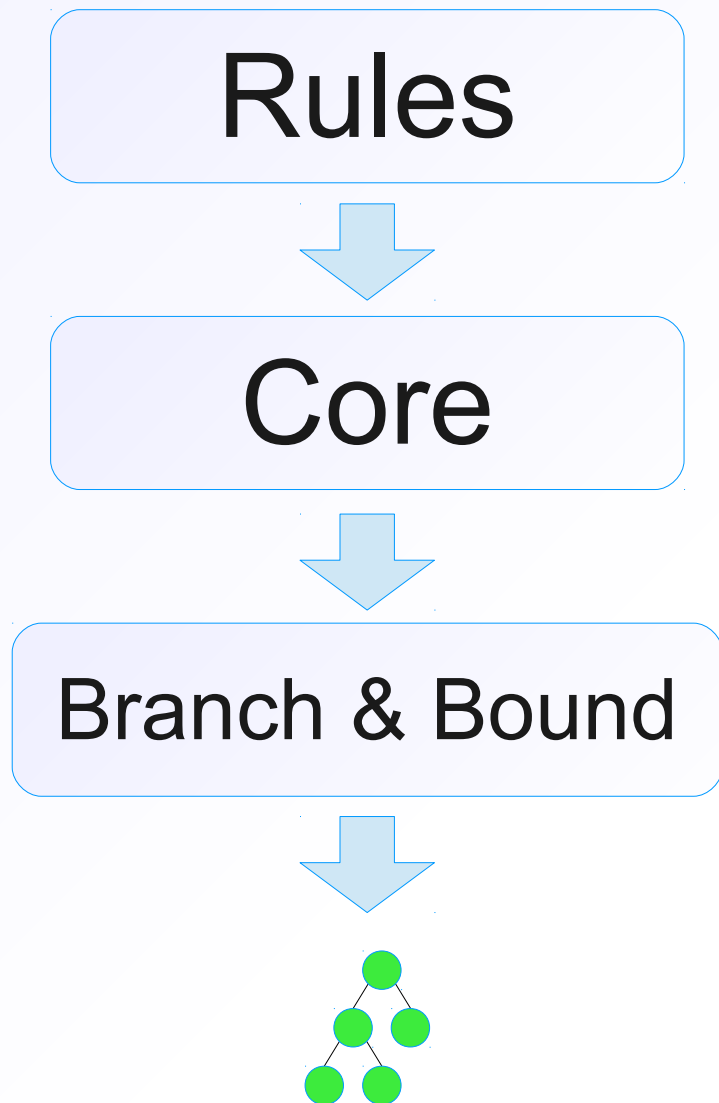


- But the algorithms produce trees with back edges
Turned out it is not very useful for most tasks

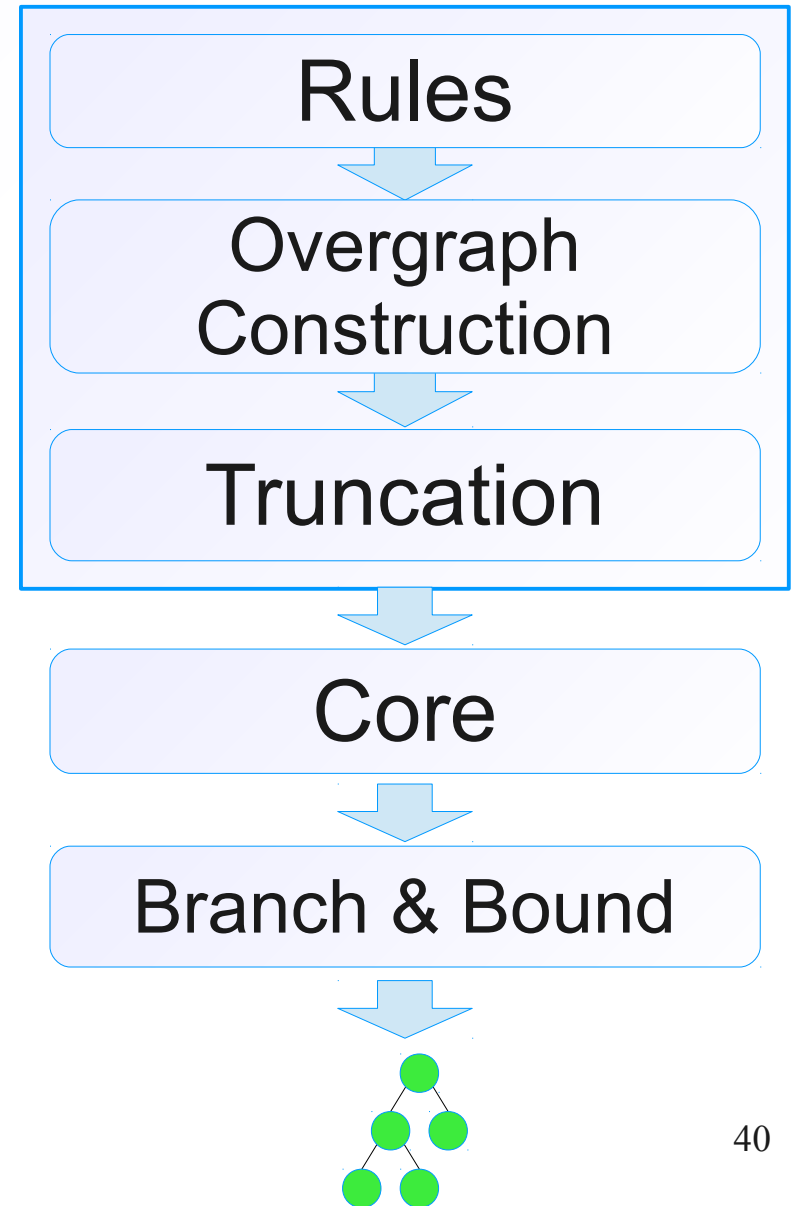
Example: Counter Systems

- The task is to find the minimal proof of a counter system's safety
- A proof is a **graph**, not a tree with back edges
- MRSC uses **cross edges** to simulate graphs
- But overgraphs may be still useful because they enable **truncation**

Experiment with Counter Systems

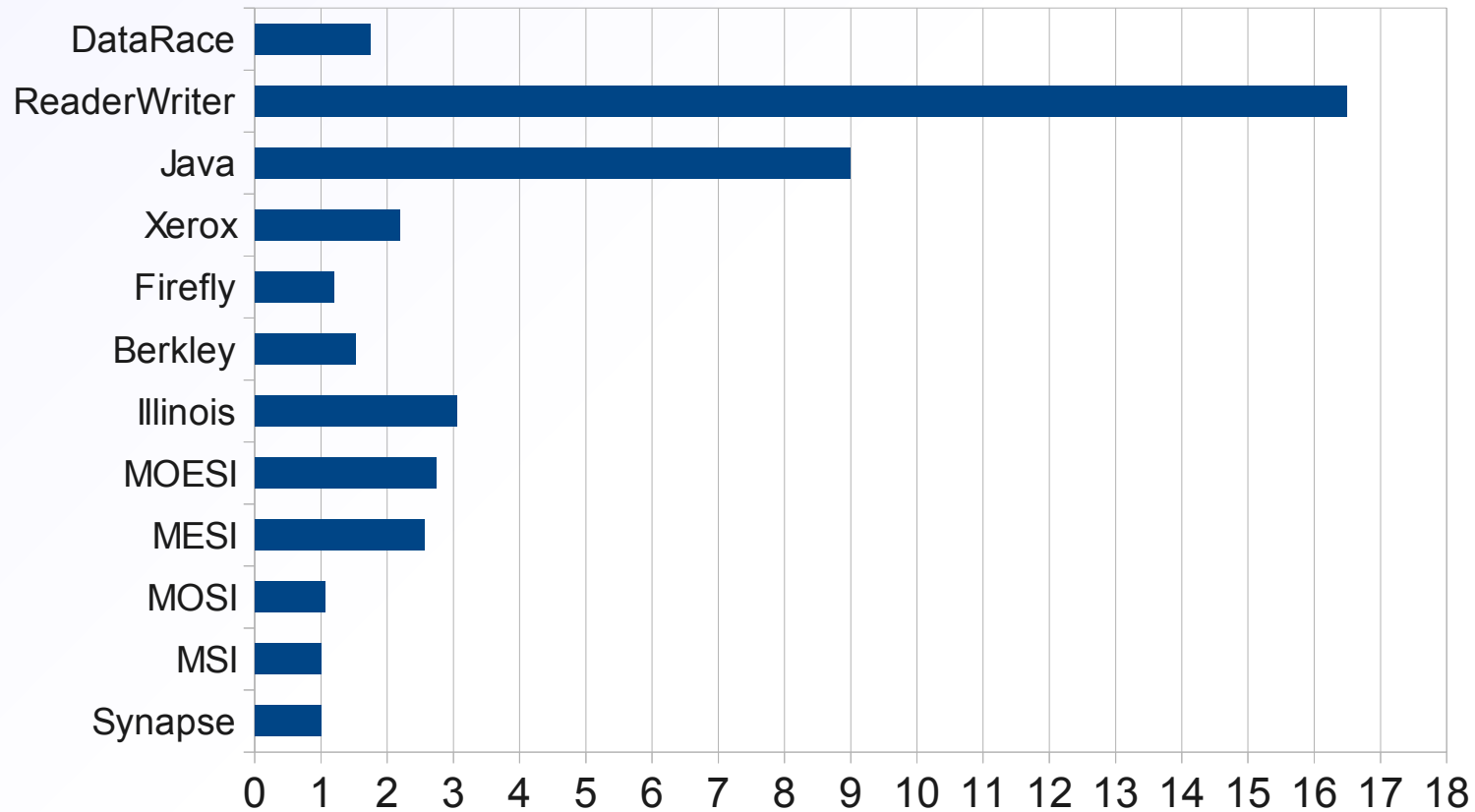


VS



Experimental Results

Improvement (times)



(in terms of the number of visited nodes)

Why overgraphs were useful?

- We could compute **sets of successors**
- We could **truncate** an overgraph

An overgraph contains a lot of **information** about relations between configurations

This is even more important than its compactness

Further Work

- Experiments with subgraph-producing residualization algorithms
 - need graph-based language
 - tree-producing algorithm seems unsuitable for real-world tasks
- Searching for heuristics (whistles etc) useful for overgraph representation
- Applying overgraphs to higher-level supercompilation

Conclusions

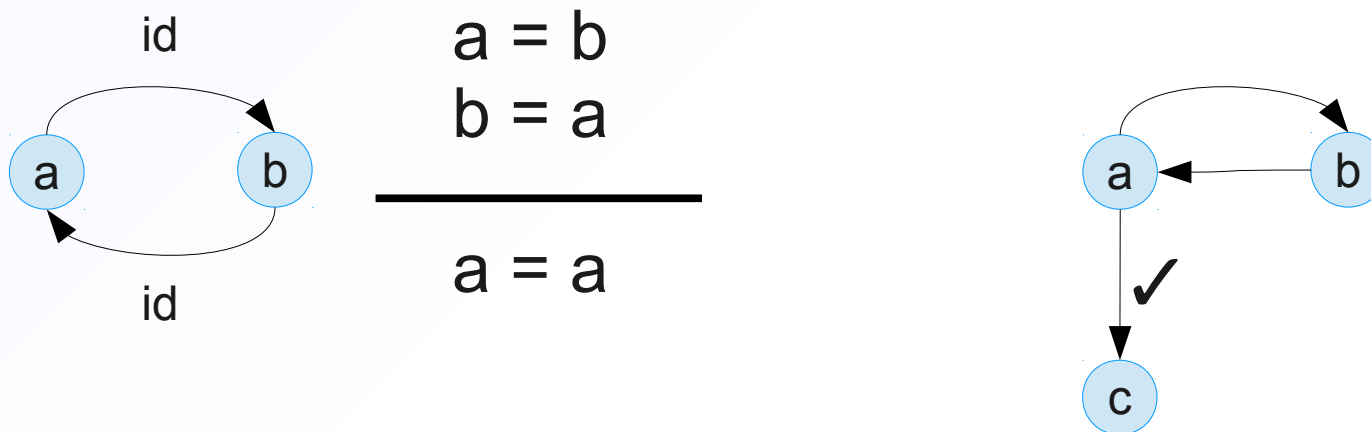
We suggested the Overgraph representation

- An Overgraph is a very compact representation
- Rules, Whistles and Residualization were generalized to Overgraphs
- The implementation has shown its usefulness
 - Caching residualization algorithm
 - Truncation for counter systems
- Overgraph contains a lot of information, so it is possible to analyze multiple graphs at once

Please return to the previous slide

Correctness

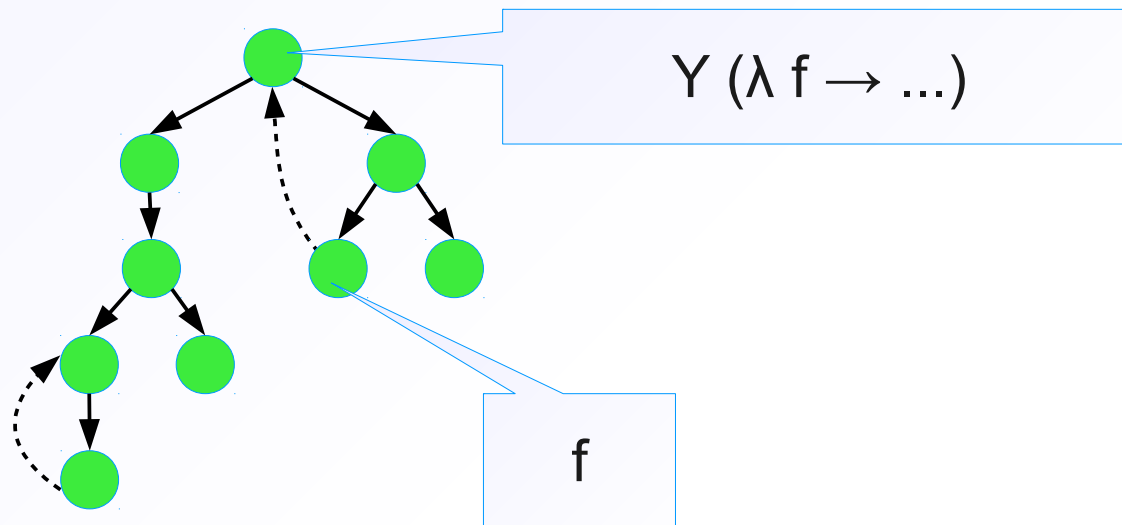
- It is possible that not all of the trees extracted from an overgraph represent correct programs



- Usually it is not a problem for single-level supercompilation

Language used in experiments

- The language is essentially based on trees with back edges



- Higher order
- Explicit fixed point combinator
- No let-expressions

Overgraph vs E-PEG

- Essentially the same idea applied to different domains
- We work with functional languages, so we have a clear recursion rather than incomprehensible cycles
- We don't have symmetric equalities
- We decided to residualize to trees, they naturally “residualize” to graphs
 - Should we do the same?

There are no more slides